



Model Compression

Present by Chun-Tse Hsiao

2007/05/02

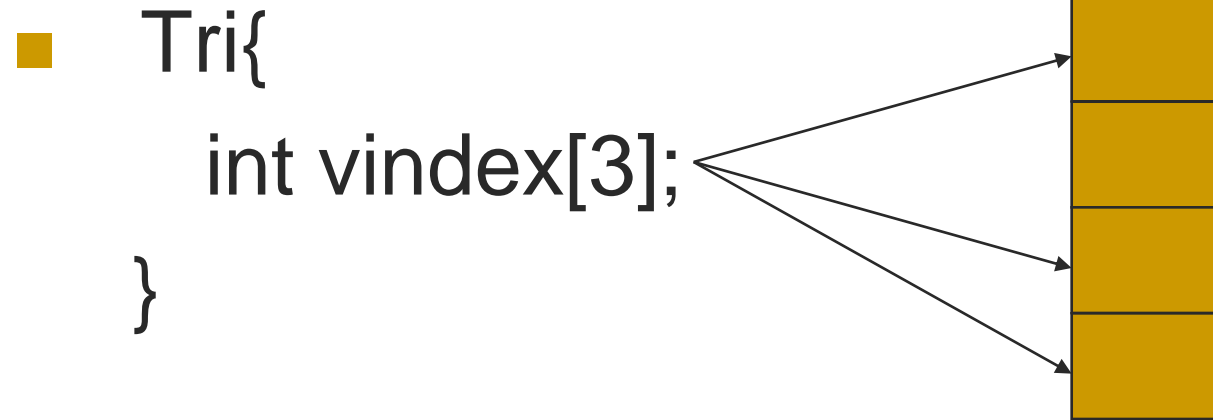
[Model Data]

- Geometry data.
- Connectivity data (Topology).
- Property data.

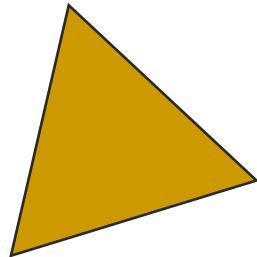
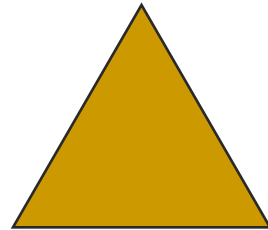
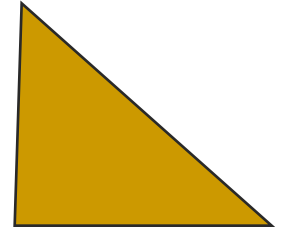
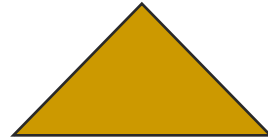
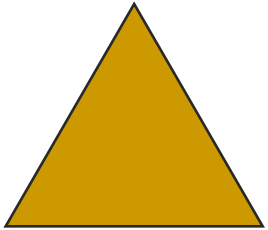
- In this presentation, we discuss compression of topology data.

[Ordinary model representation]

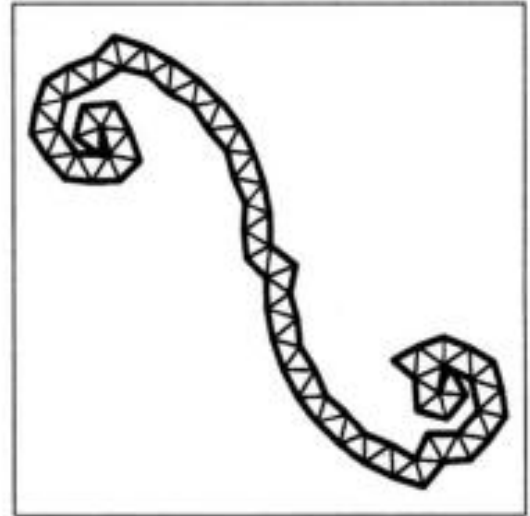
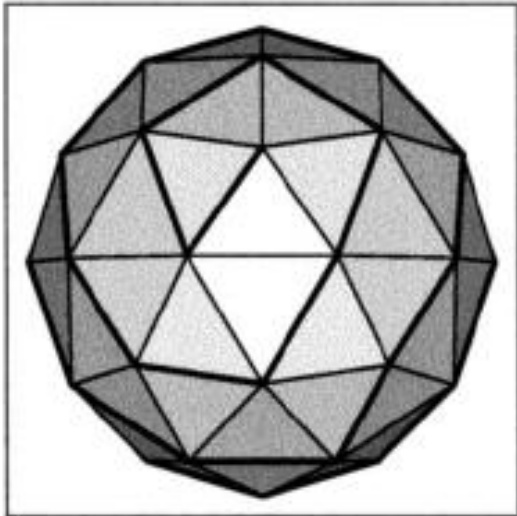
Vertex position array



- The representation above can represent arbitrary triangle model.



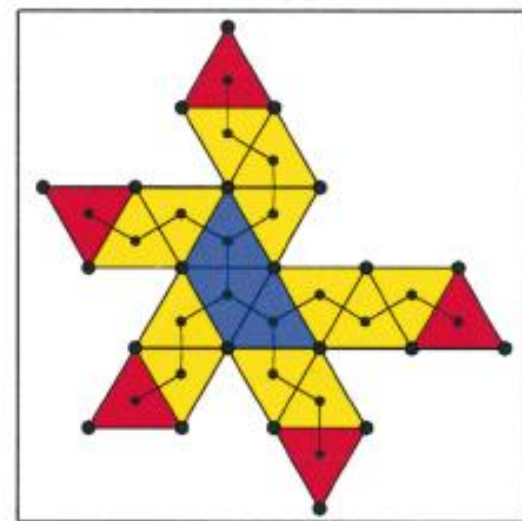
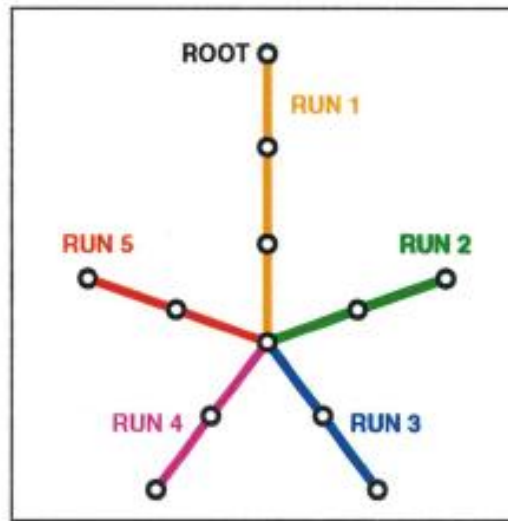
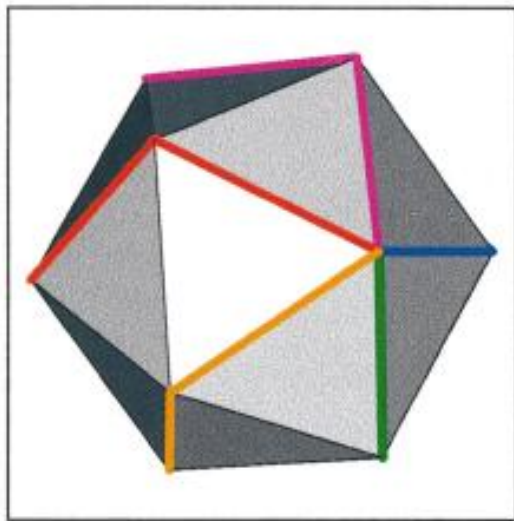
[Triangle strips]



[Complexity of topology]

- Topology complexity is proportional to total edge number.
- Edges describe Topology information.

[Previous work]




Topology can be represented by vertex spanning tree and face spanning tree! Total edge number of two tree is equal to total edge number of original model.

[Edgebreaker

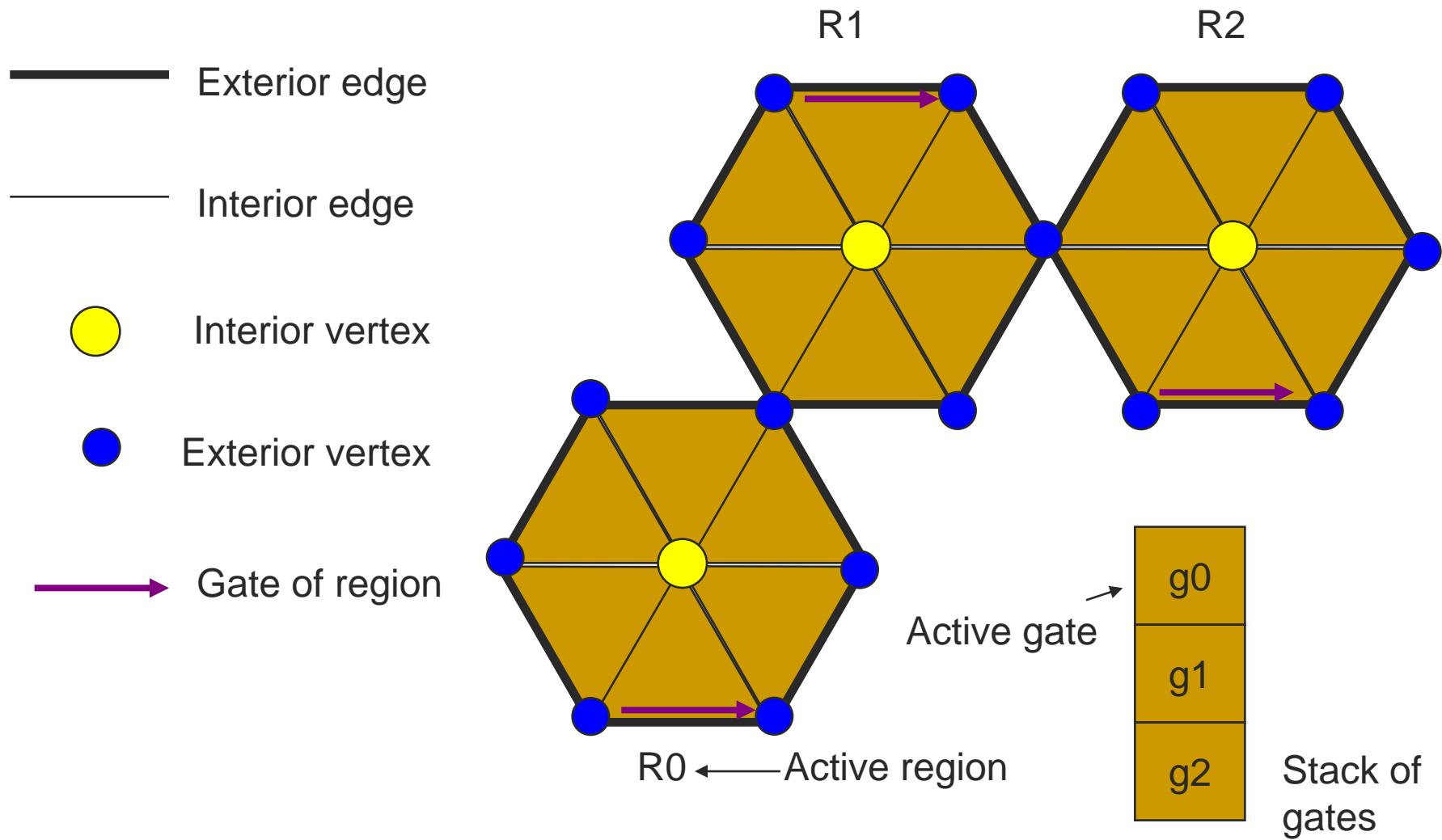
]


[Definition]

- Simple mesh
 - 2-manifold triangle mesh
 - Connected & Orientable
 - Have no handle
 - Have no boundary or have a connected, manifold, close curve boundary.

- 
- A large black left square bracket and a large yellow right square bracket are positioned at the top of the slide, with a thin yellow horizontal line extending between them across the width of the content area.
- Edgebreaker compression algorithm performs a **series of steps**.
 - Each steps remove one triangle from current mesh.
 - The remaining portion mesh is composed of one or several simple mesh region.

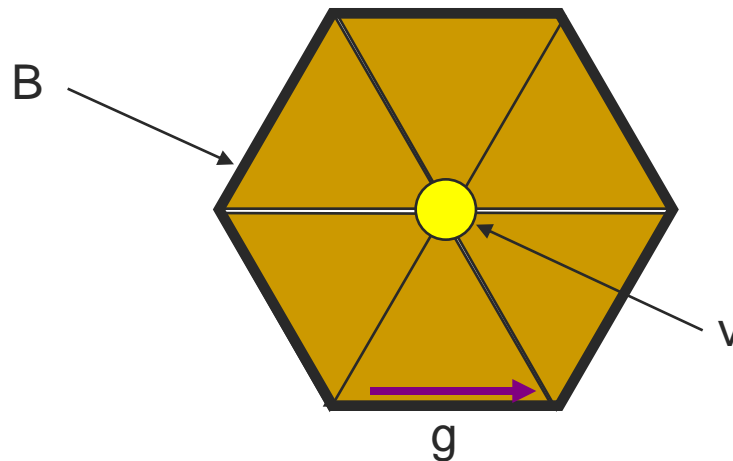
[3 regions example]



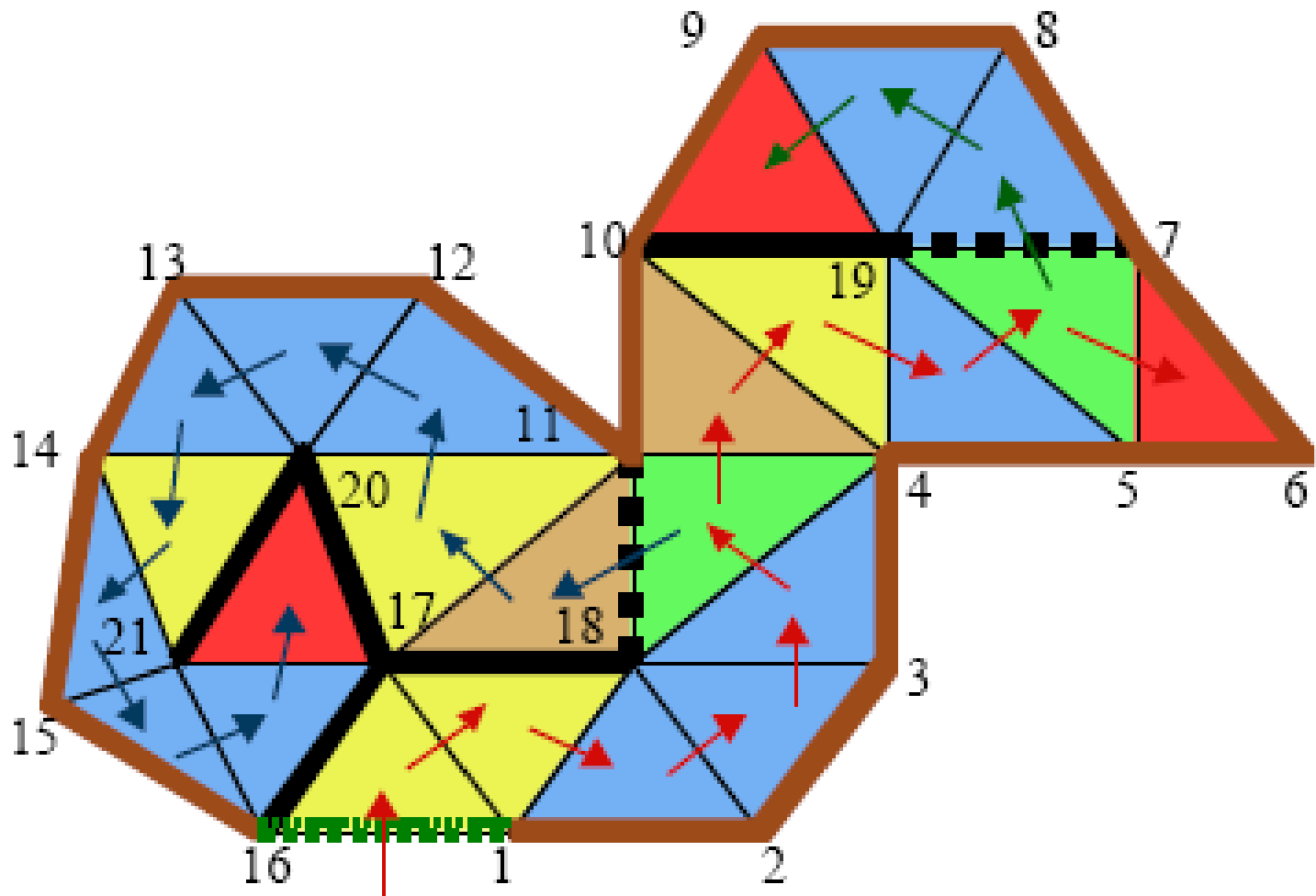
- 
- A large black left square bracket is on the left side of the slide, and a large yellow right square bracket is on the right side. A horizontal line with a light green gradient runs across the slide, starting from the left bracket and ending at the right bracket.
- The regions processed in the same order as gates in the gate stack.
 - R0 will compress first, R1 next, and finally R2.
 - **Each step will remove one triangle from active region.** It may introduce new region! New region will be tracked by gate stack.

■ Notations

- Border of active region : B
- Gate of active region : g
- Vertex in the same triangle not bound g : v

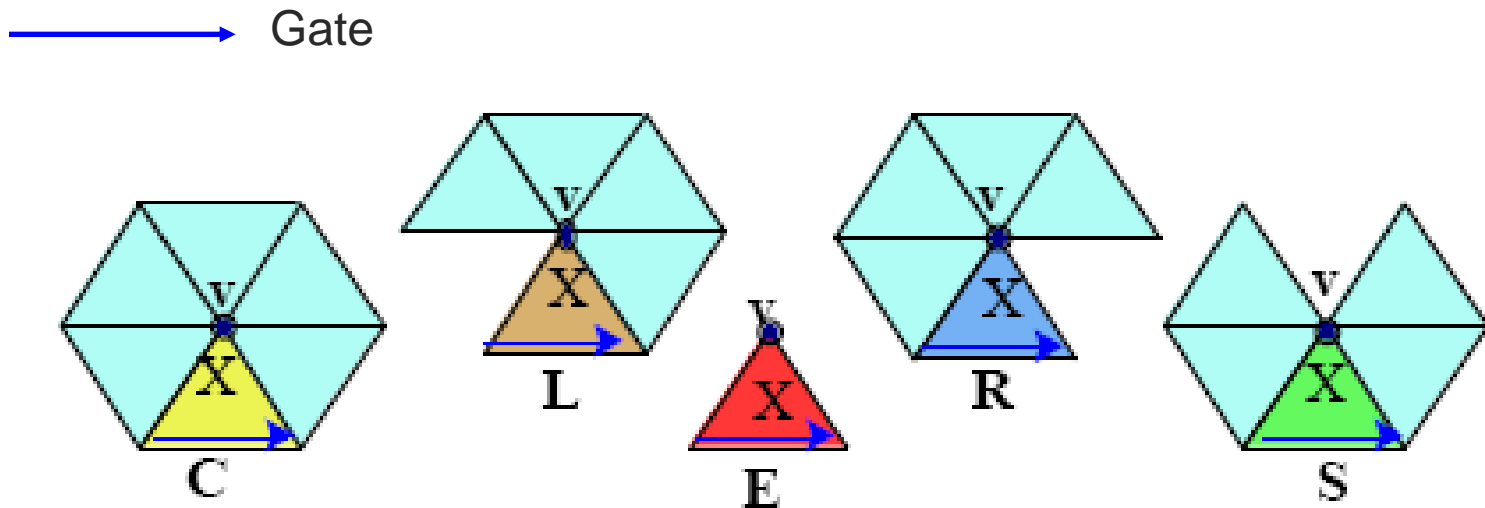



[Compressing simple meshes]



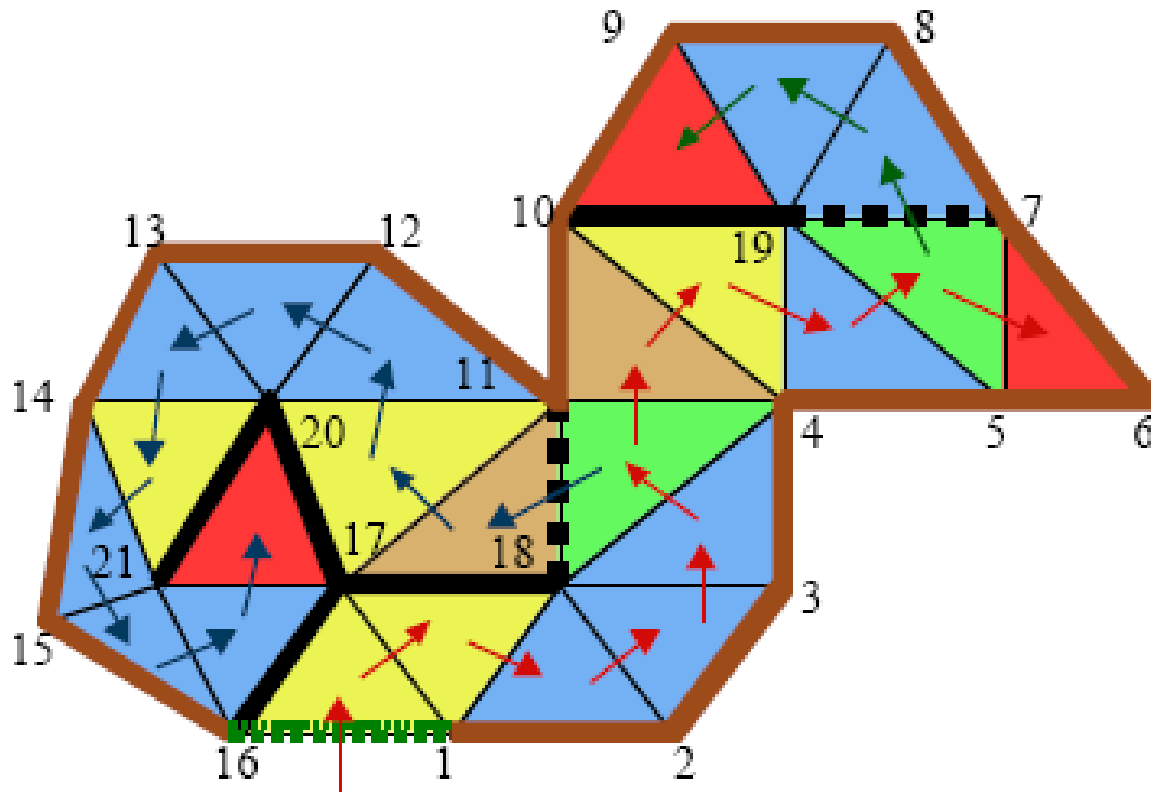
[Five operation of edge breaker.]

- Depend on the relation of v , g , B ...there will be five possible operations.



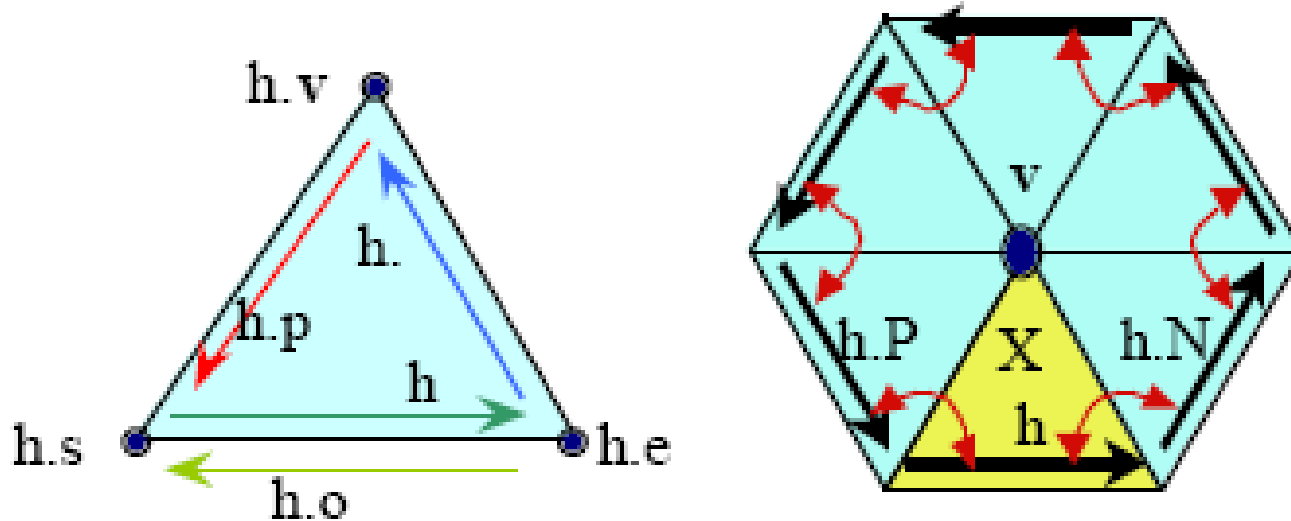
- 
- A large black left square bracket and a large yellow right square bracket are positioned at the top of the slide, with a horizontal line between them. The line is light green on the left and transitions to light yellow on the right.
- We record series of operations, denote it by H .
 - We record series of vertices, in the order in which they are reached by C operations, denote it by P .
 - If mesh has boundary, P is initialized to references of vertices of initial loop B .
 - H actually represent connectivity information of the mesh!

[Compressing simple meshes]



H = CCRRRSLCRSERREL CRRRCRRRE

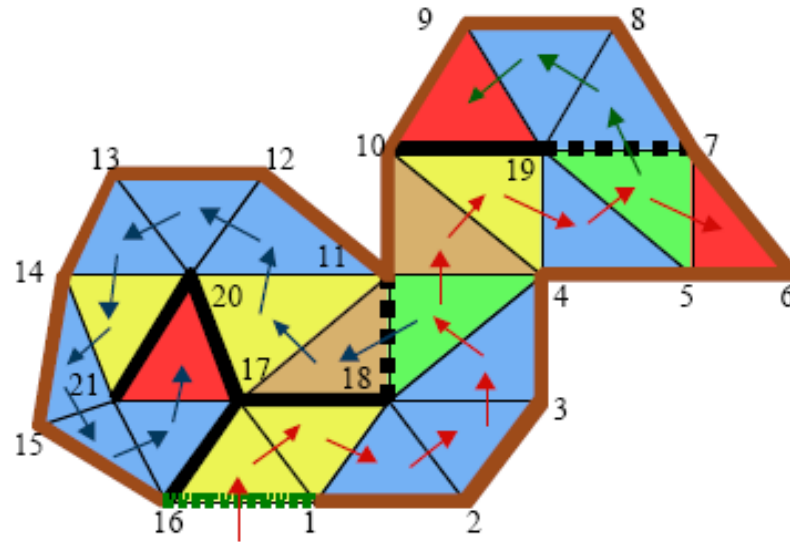
[Half-edge data structure]



Each vertex v has flag $v.m = \text{true}$ if it's visited before.

Each half-edge h has flag $v.m = \text{true}$ if it's in bounding loop of remaining portion of mesh

[Initialize]



$P = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16\}$

$H = \{\}$

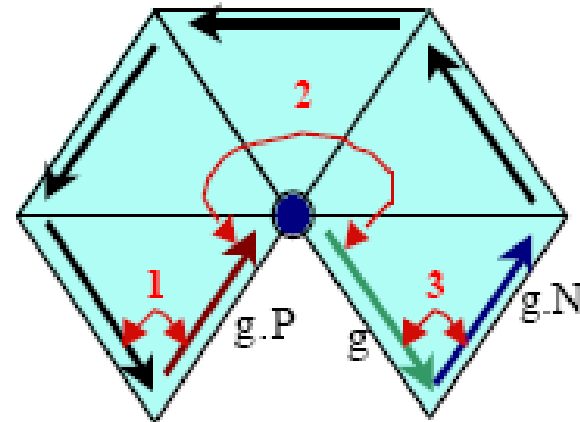
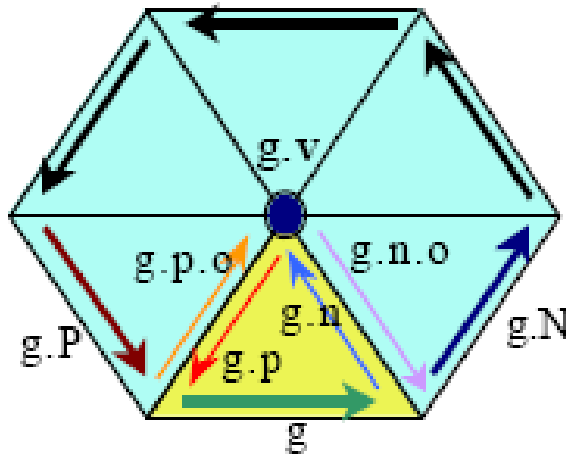
Set $V_1 \sim V_{16}$'s visit flag to true.

Set boundary flag of half-edges on boundary to true.

$g = (16, 1)$

Stack = {g}

Case C



```

H=H|C;
P=P|g.v;
g.m=0; g.p.o.m=1;
g.n.o.m=1; g.v.m=1;
g.p.o.P=g.P; g.P.N=g.p.o;
g.p.o.N=g.n.o; g.n.o.P=g.p.o;
g.n.o.N=g.N; g.N.P=g.n.o
g=g.n.o; StackTop=g;

```

```

# append C to history
# append v to P
# update flags

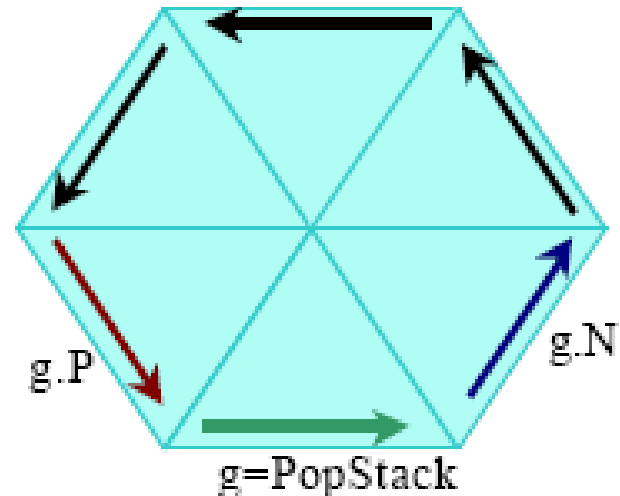
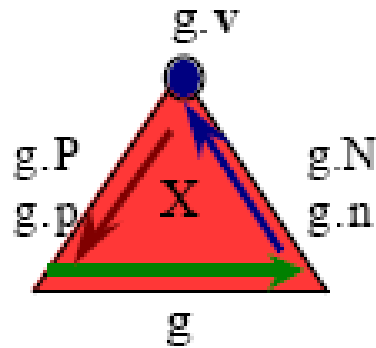
```

```

# fix red link 1 in B
# fix red link 2 in B
# fix red link 3 in B
# move gate

```

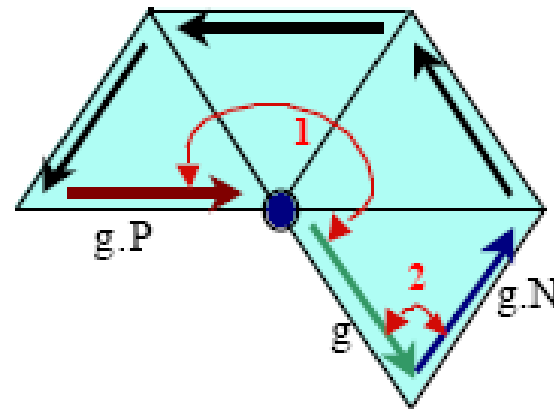
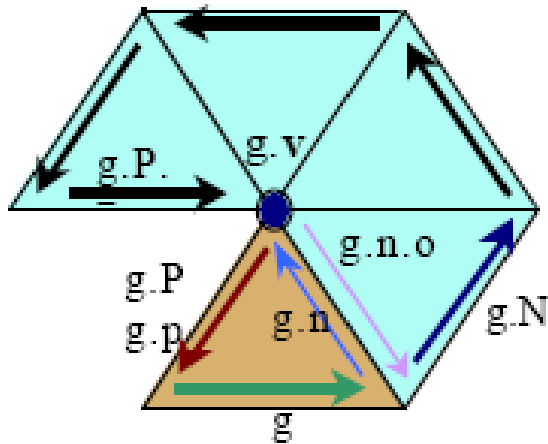
[Case E]



$H=H|E;$
 $g.m=0; g.n.m=0; g.p.m=0;$
 $PopStack; g=StackTop;$

append E to the history
unmark edges
pop stack: next region

Case L



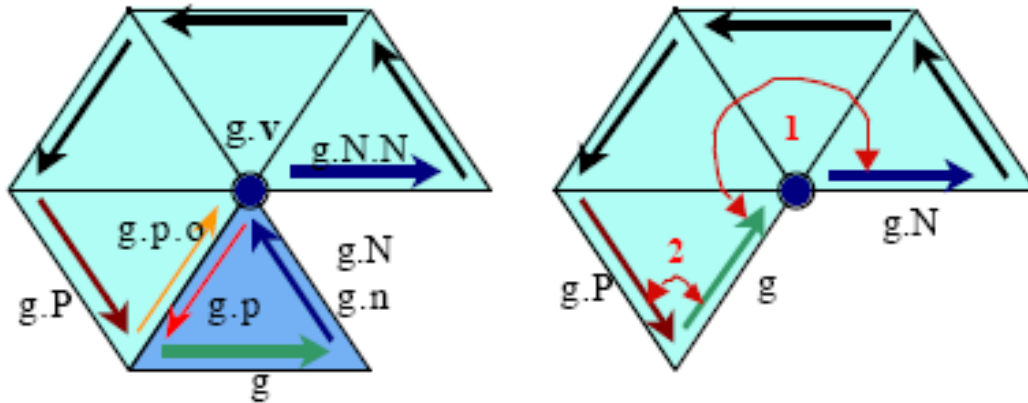
```

H=H|L;
g.m=0; g.P.m=0; g.n.o.m=1;
g.P.P.N=g.n.o; g.n.o.P=g.P.P;
g.n.o.N=g.N; g.N.P=g.n.o;
g=g.n.o; StackTop=g;
    
```

```

# append L to history
# update marks
# fix red link 1 in B
# fix red link 2 in B
# move gate
    
```

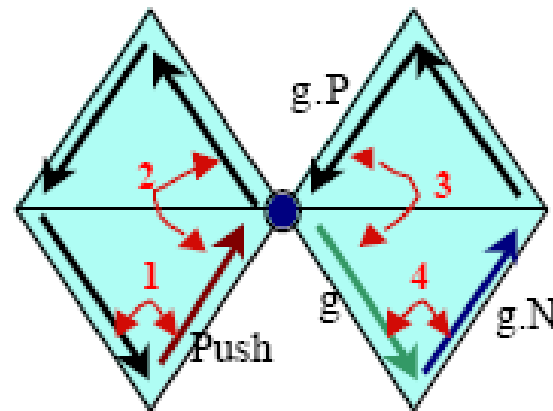
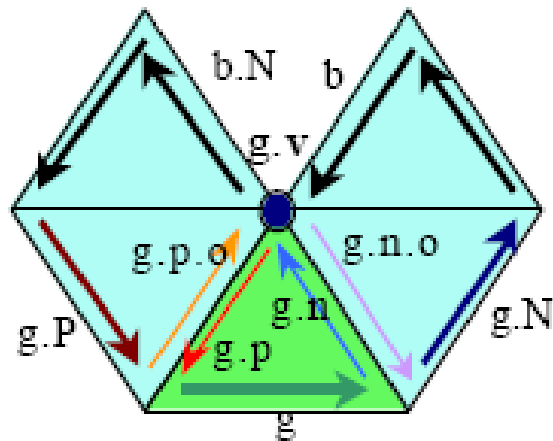
Case R



$H = H \cup R;$
 $g.m = 0; g.N.m = 0; g.p.o.m = 1;$
 $g.N.N.P = g.p.o; g.p.o.N = g.N.N;$
 $g.p.o.P = g.P; g.P.N = g.p.o;$
 $g = g.p.o; StackTop = g;$

append R to history
 # update marks
 # fix red link 1 in B
 # fix red link 2 in B
 # move g

Case S



```

H=H|S;
g.m=0; g.p.o.m=1; g.n.o.m=1;
b=g.n;
WHILE NOT b.m DO b=b.o.p;
g.P.N=g.p.o; g.p.o.P=g.P;
g.p.o.N=b.N; b.N.P=g.p.o;
b.N=g.n.o; g.n.o.P=b;
g.n.o.N=g.N; g.N.P=g.n.o;
StackTop=g.p.o; PushStack;
g=g.n.o; StackTop=g
    
```

```

# append S to history
# update marks
# initial candidate for b
# turn around v to marked b
# fix red link 1 in B
# fix red link 2 in B
# fix red link 3 in B
# fix red link 4 in B
# push g.p.o on stack
# move g
    
```


[Decompression]

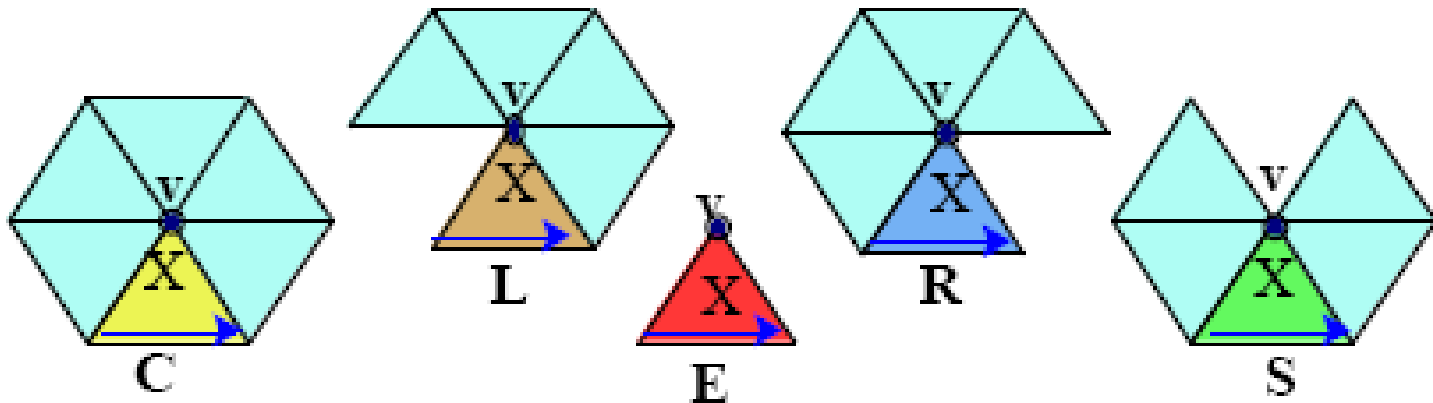
- Two pass decompression
- Preprocessing (Parse H)
 - Preprocessing $|T|$, $|VE|$, $|VI|$, and offset of all S operations in offset table O.
- Generation
 - Create triangles in the order in which they were deleted by compression process.

[Preprocessing]

- t : total number of operation. (0)
- d : $|S| - |E|$, after last E, it'll be negative. (0)
- c : $|C|$. (0)
- e : $3|E| + |L| + |R| - |C| - |S|$. Final value is $|VE|$. (0)
- s : $|S|$. (0)
- Stack : save (e, s) pairs resulting from S operations. Used during E operations to compute the offset. ($\{\}$)
- O : offset table. ($\{\}$)

[Preprocessing]

- Case S : $e -= 1; s += 1; \text{push}(e, s); d += 1.$
- Case E : $e += 3; (e', s') = \text{pop}; O[s'] = e - e' - 2; d -= 1;$
- Case C : $e -= 1; c += 1;$
- Case R : $e += 1;$
- Case L : $e += 1;$

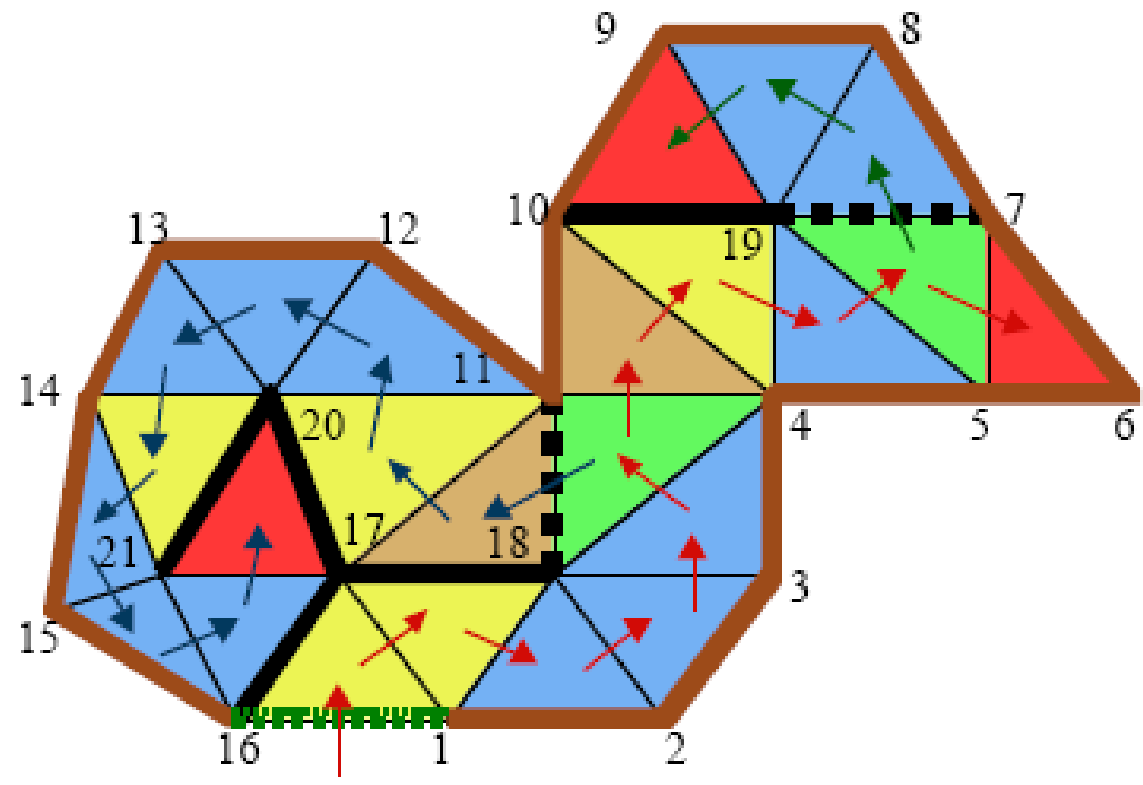


Generation

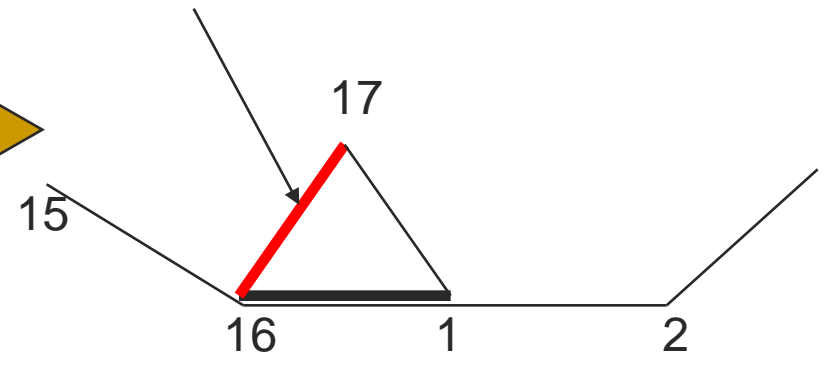
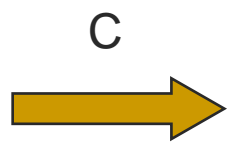
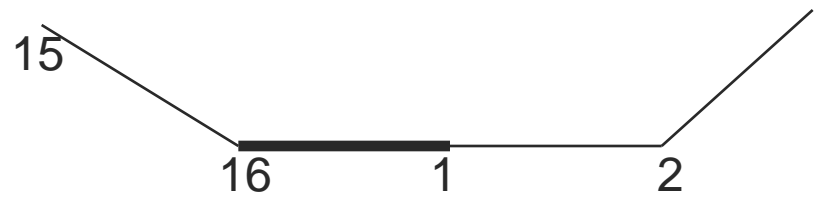
- Allocate a table of triangle-vertex incident relation TV of $|T|$ entries.
- Initialize vertex counter c to $|VE|$.
- Construct the bounding loop B of $|VE|$ vertices. (circular doubly-linked list)
- Create gate stack with single entry refer to first in B .
- Initialize t (triangle count) , s (number of S op) to zero.

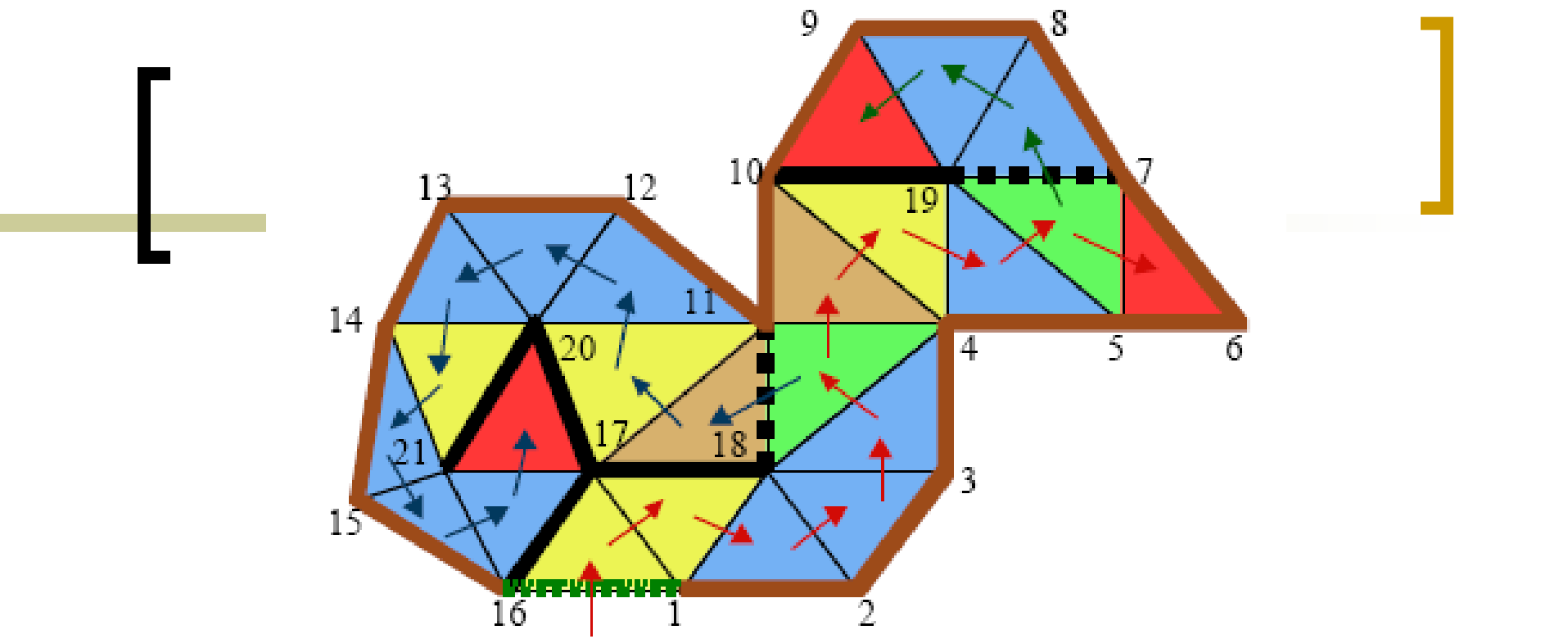
[Generation]

- Reading H .
- For each operation, increase t , store vertex index to $TV[t]$.
- Update B , g (gate), and stack if necessary.



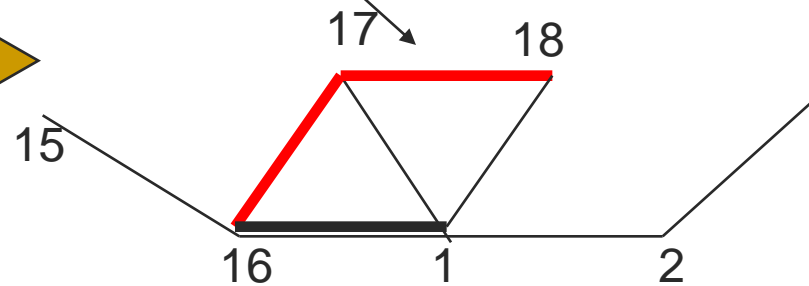
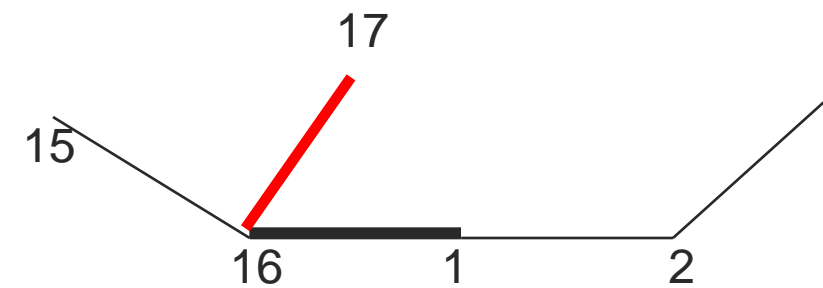
Add one edge

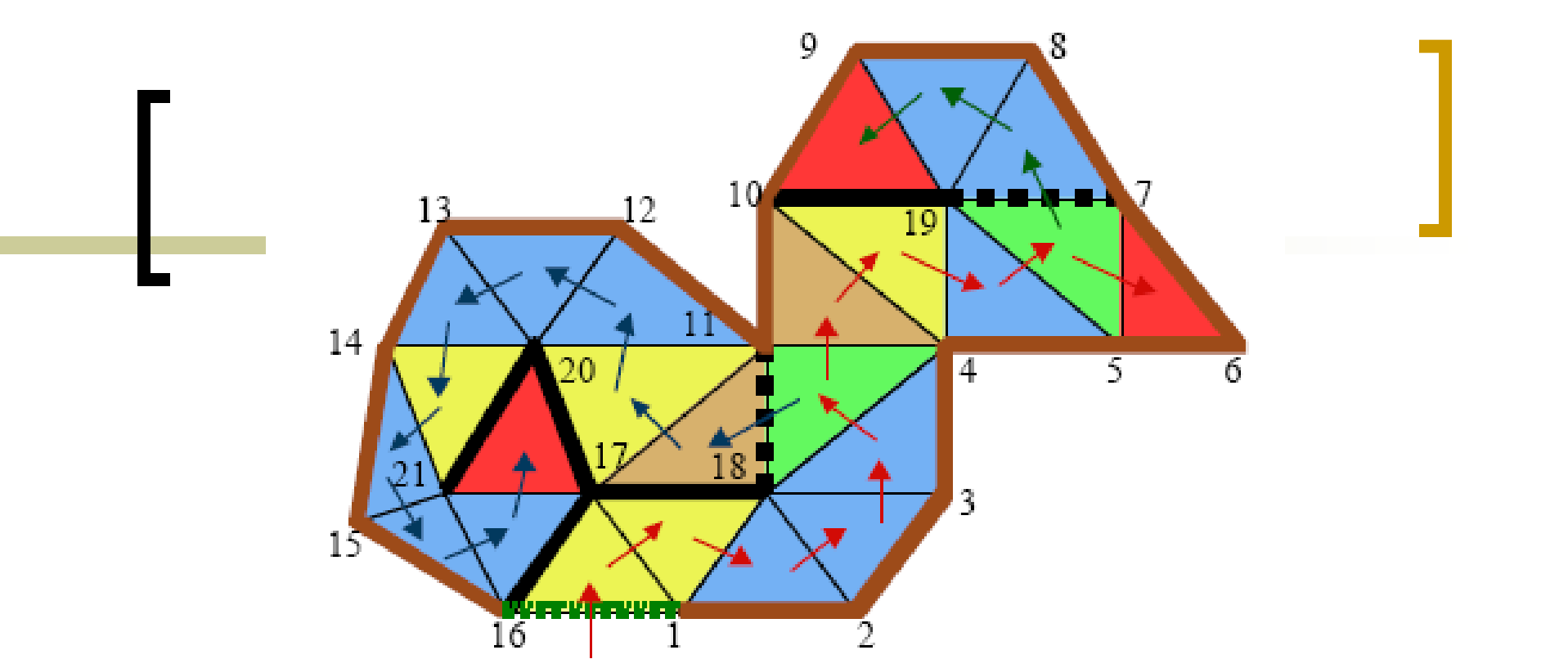




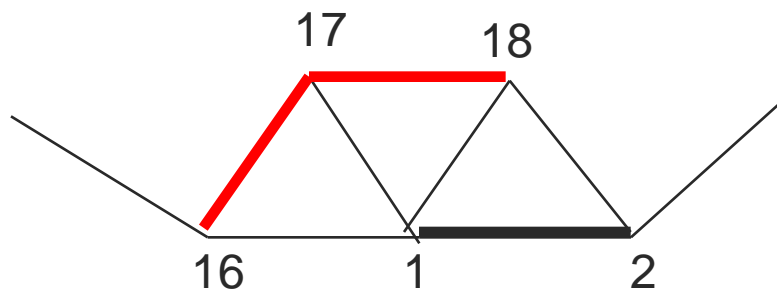
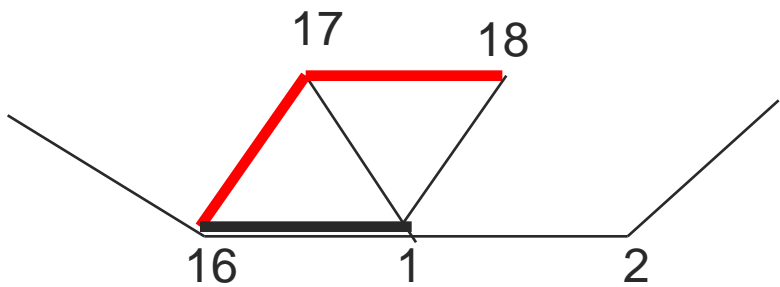
Add one edge

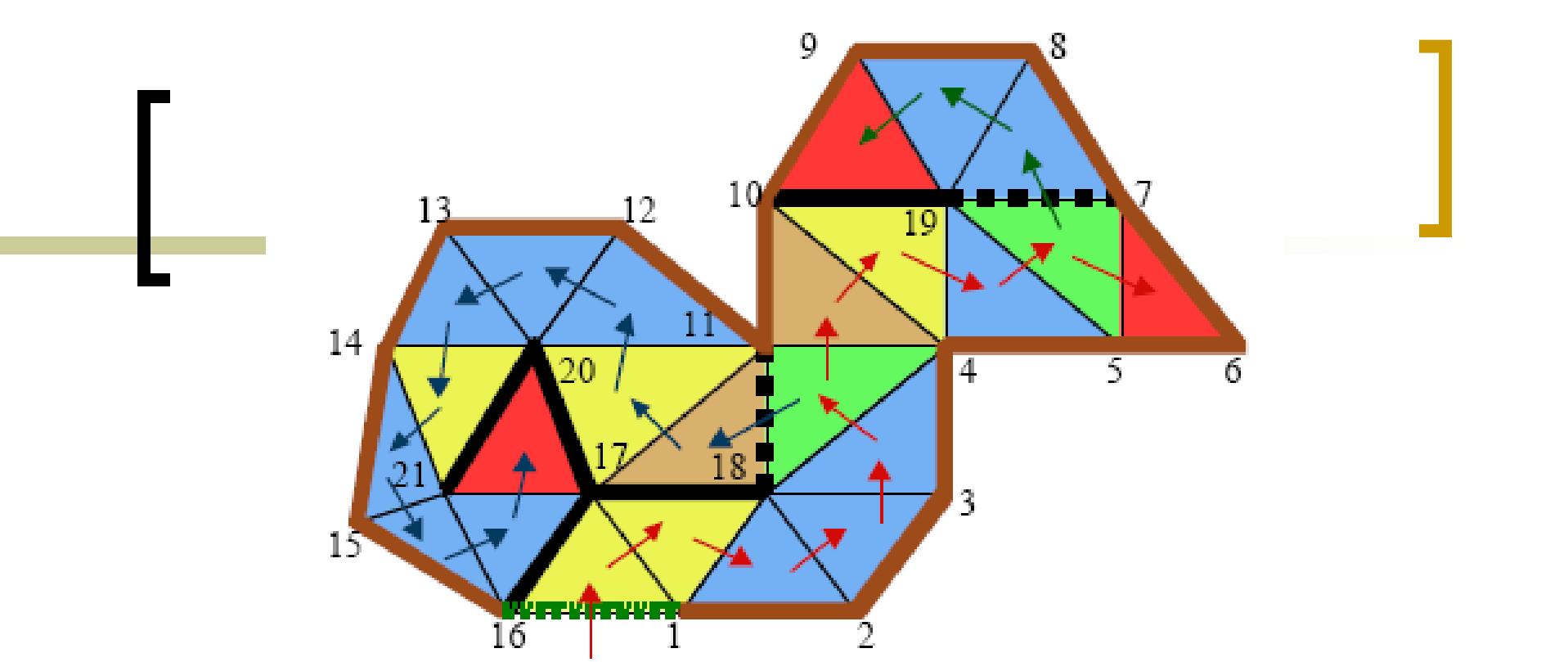
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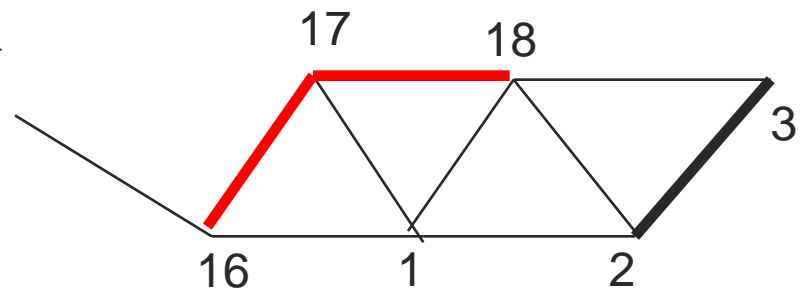
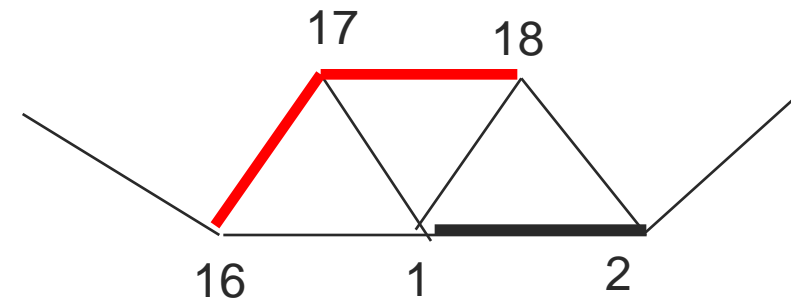


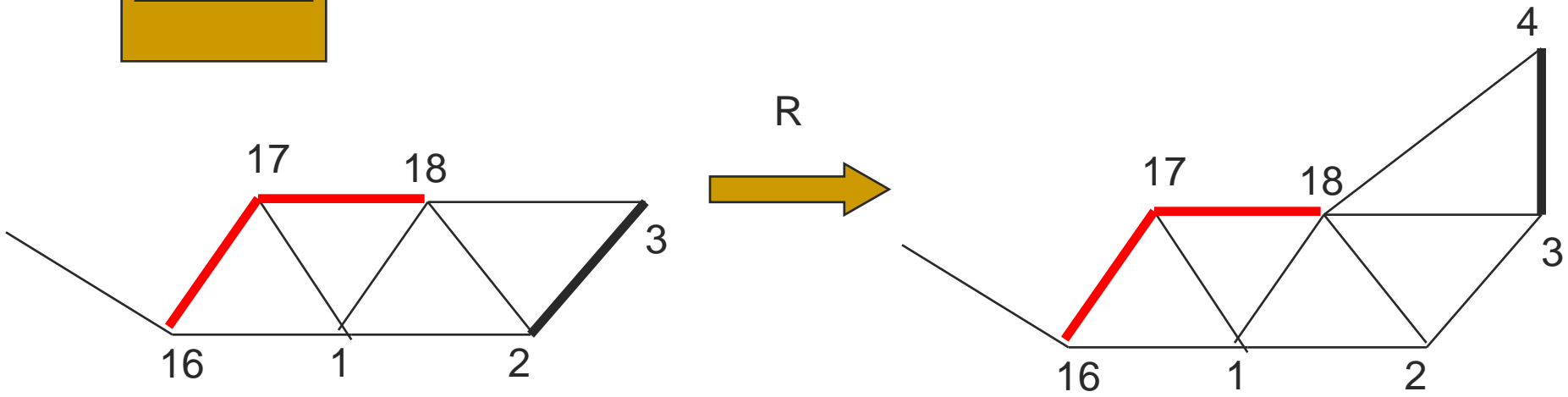
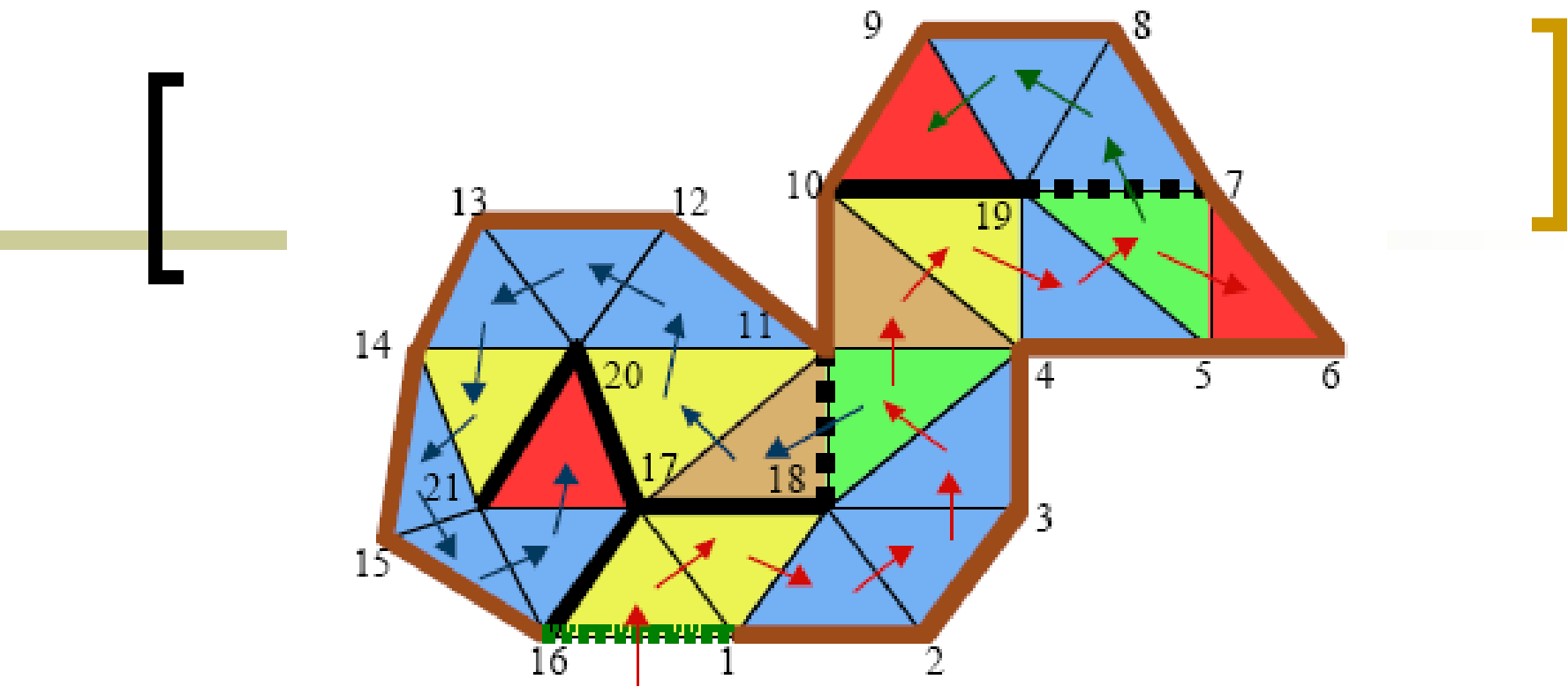
R

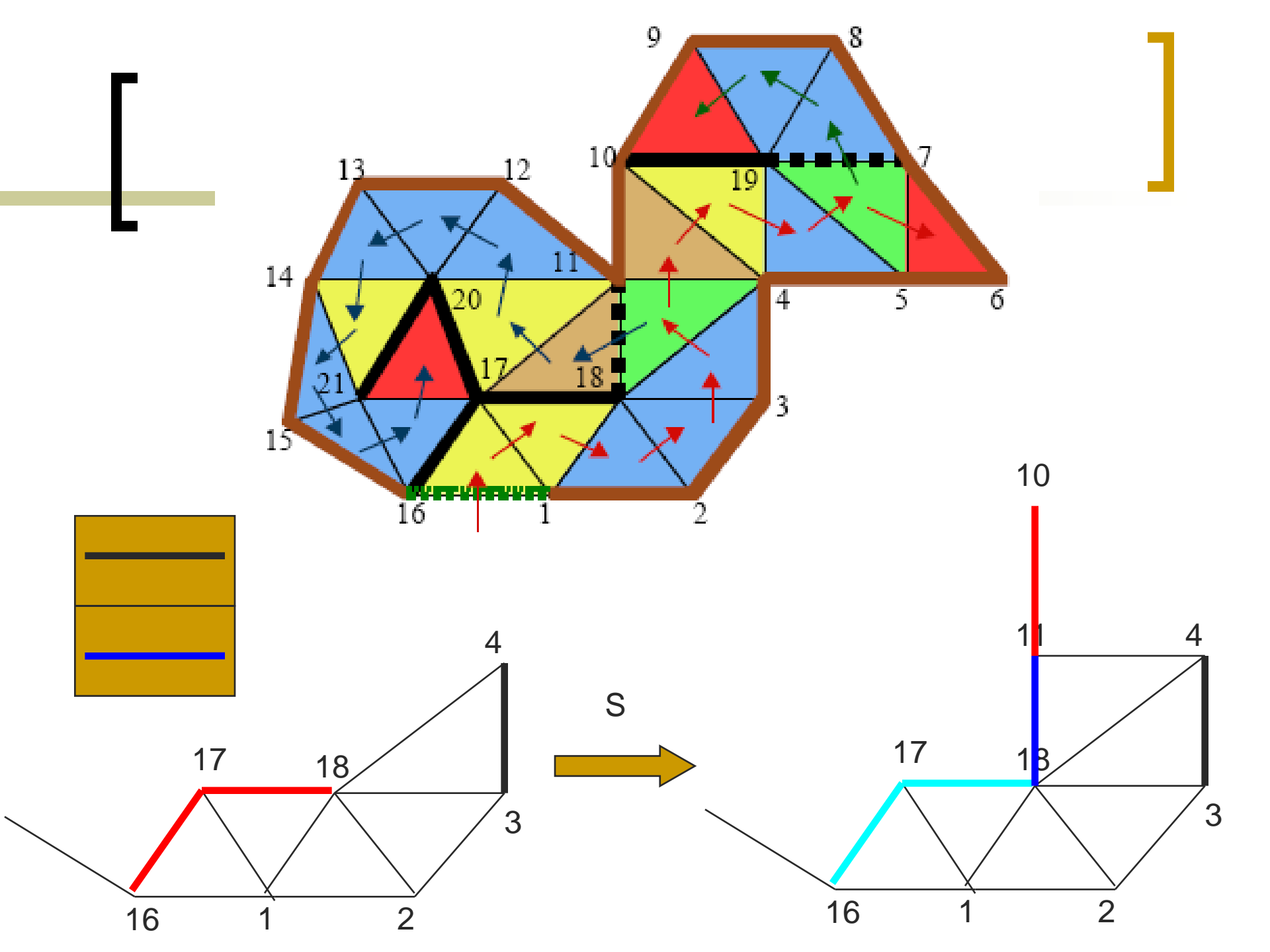


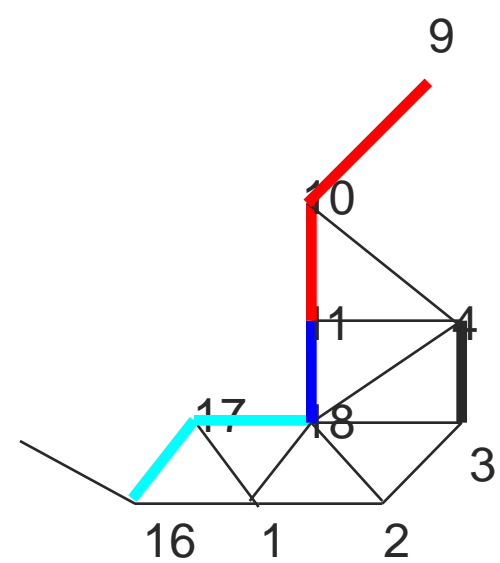
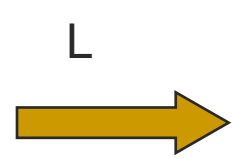
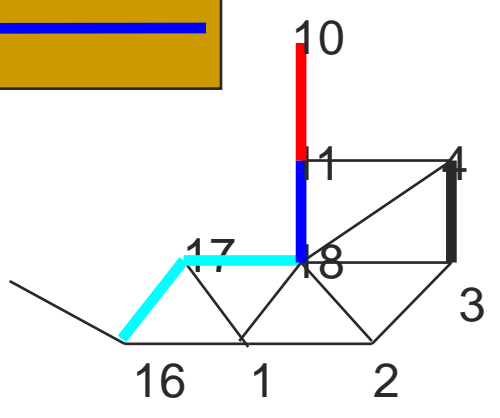
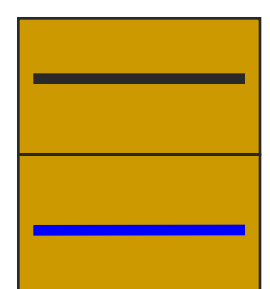
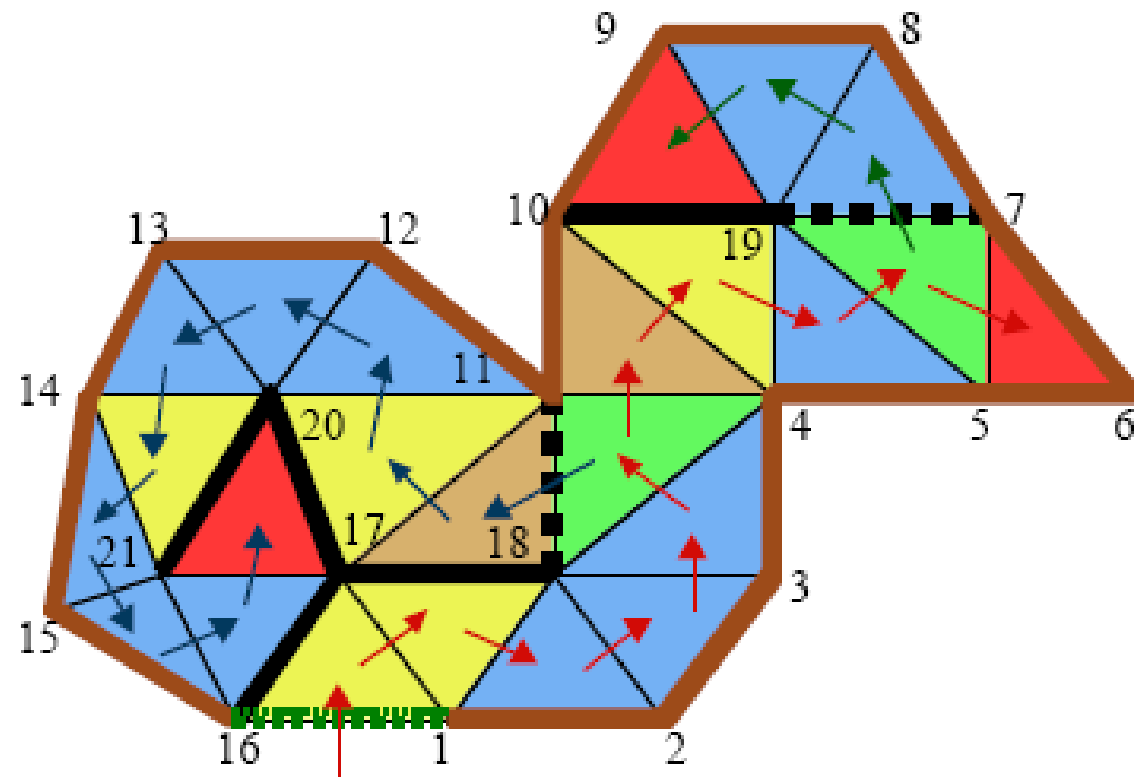
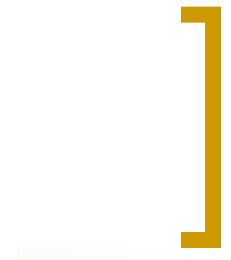
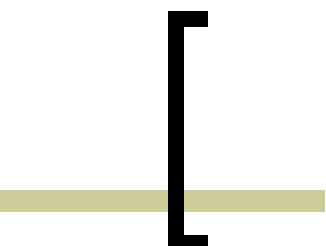


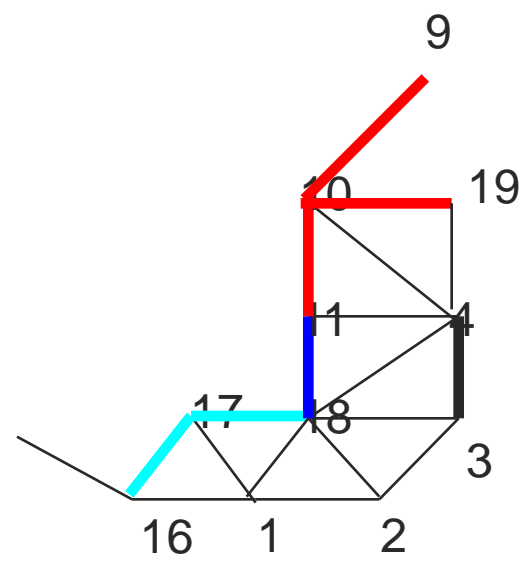
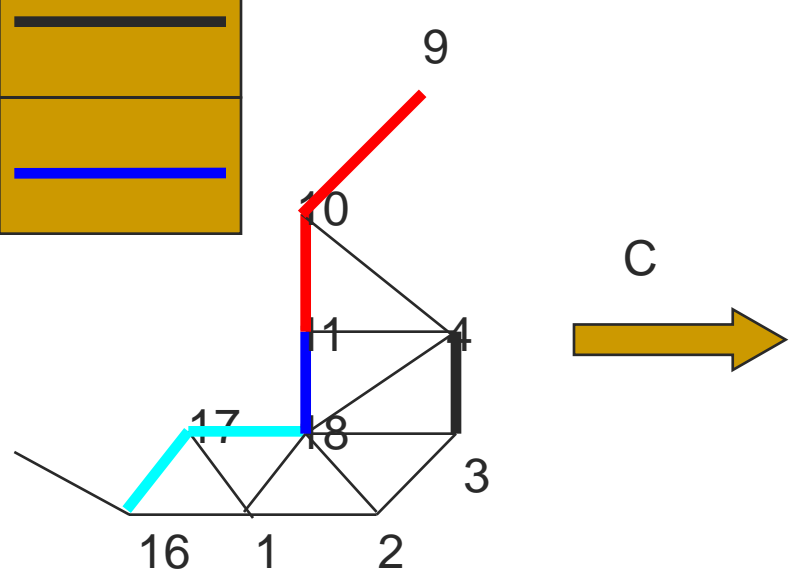
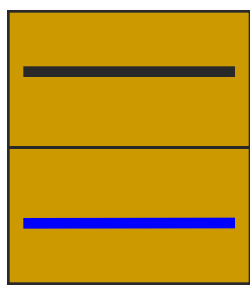
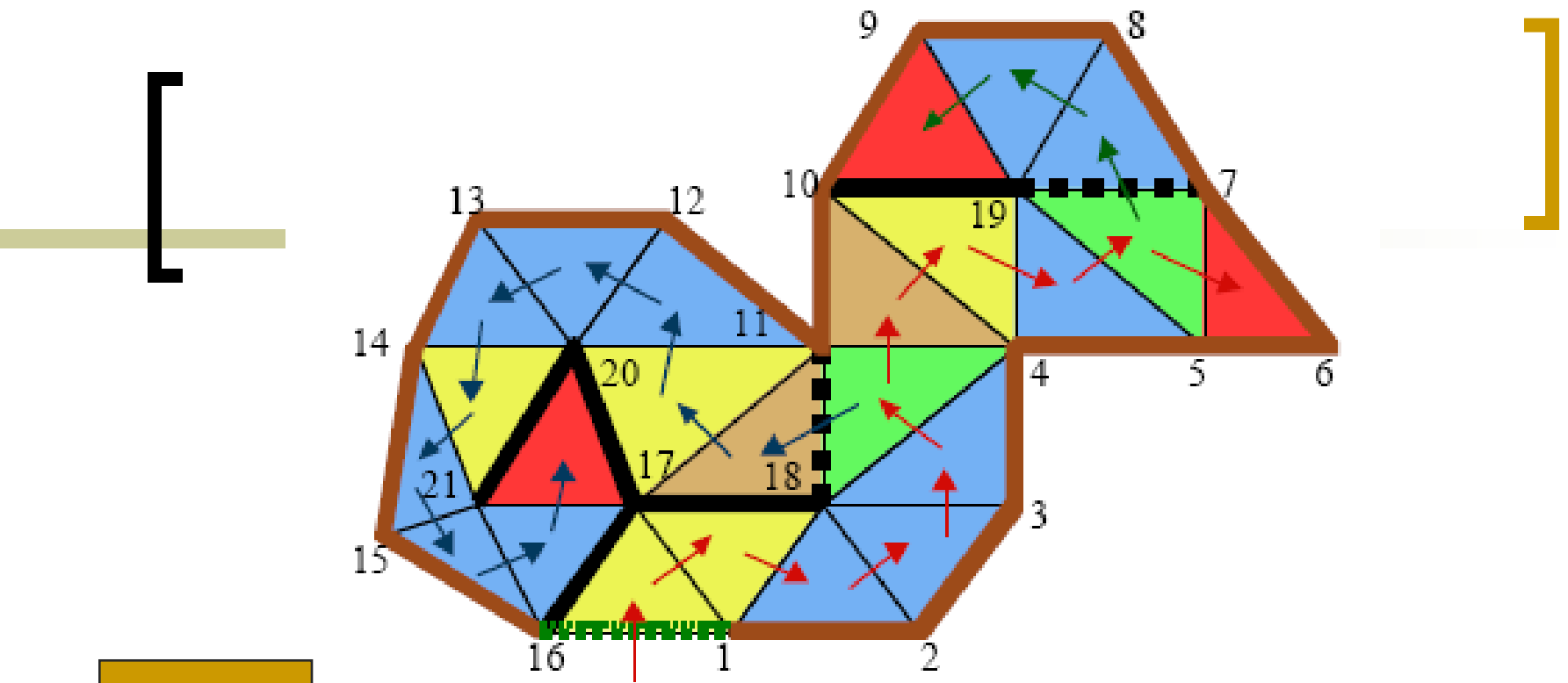
R

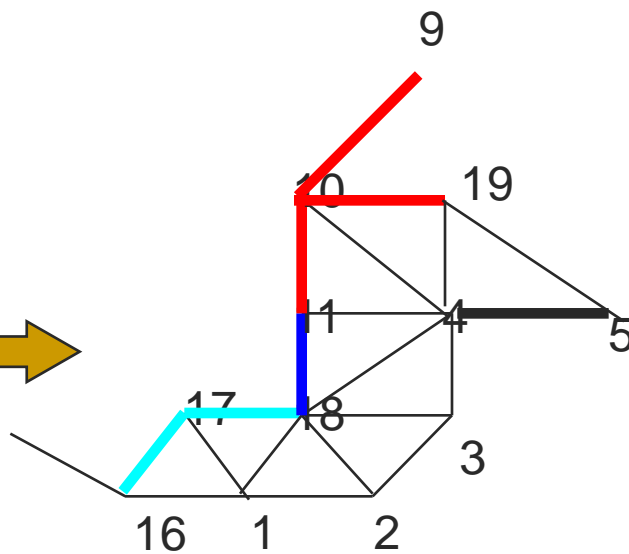
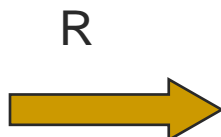
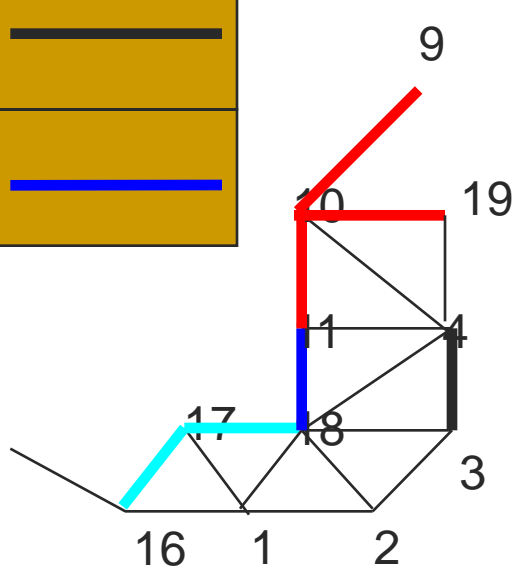
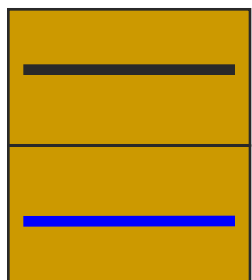
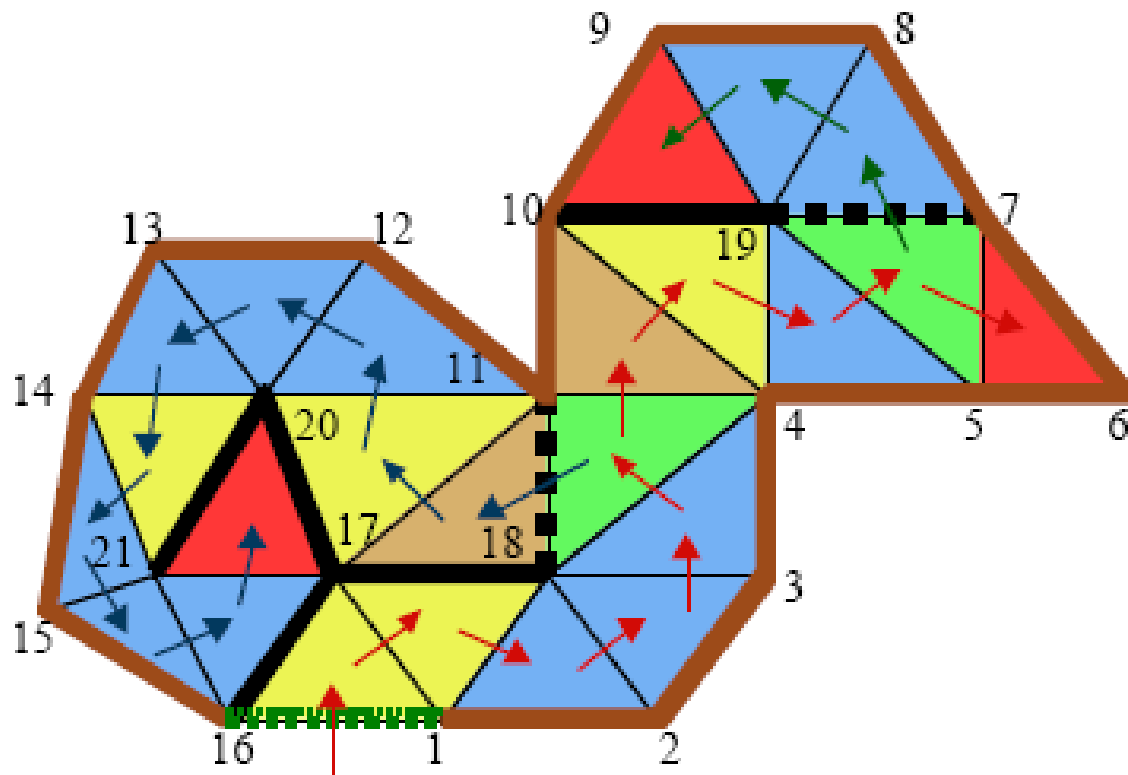


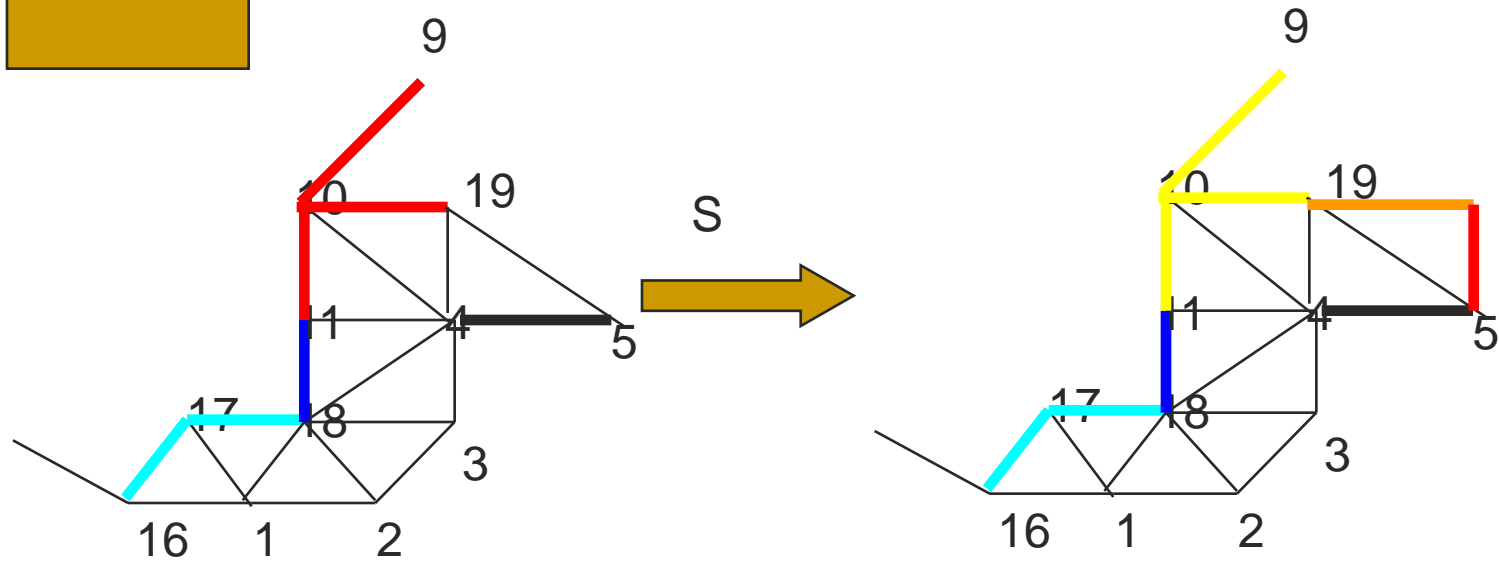
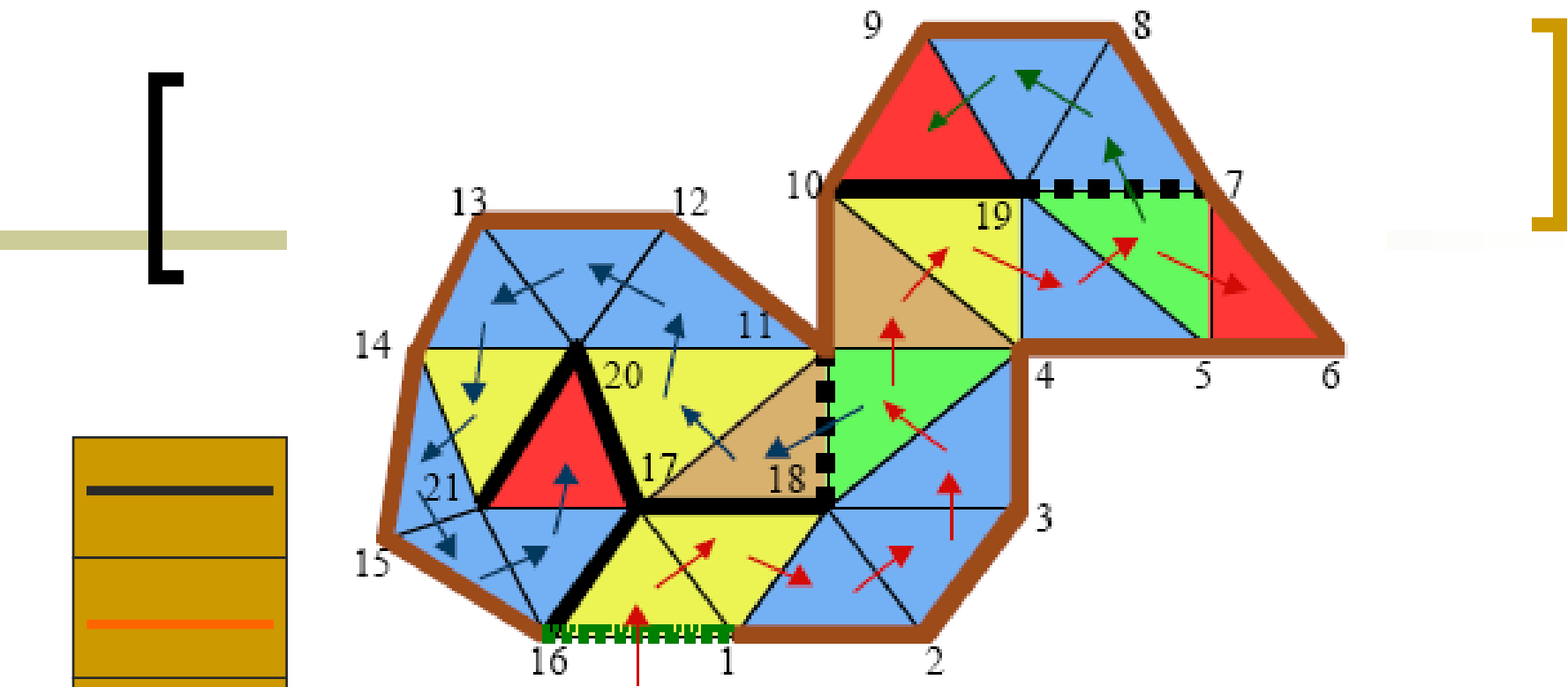


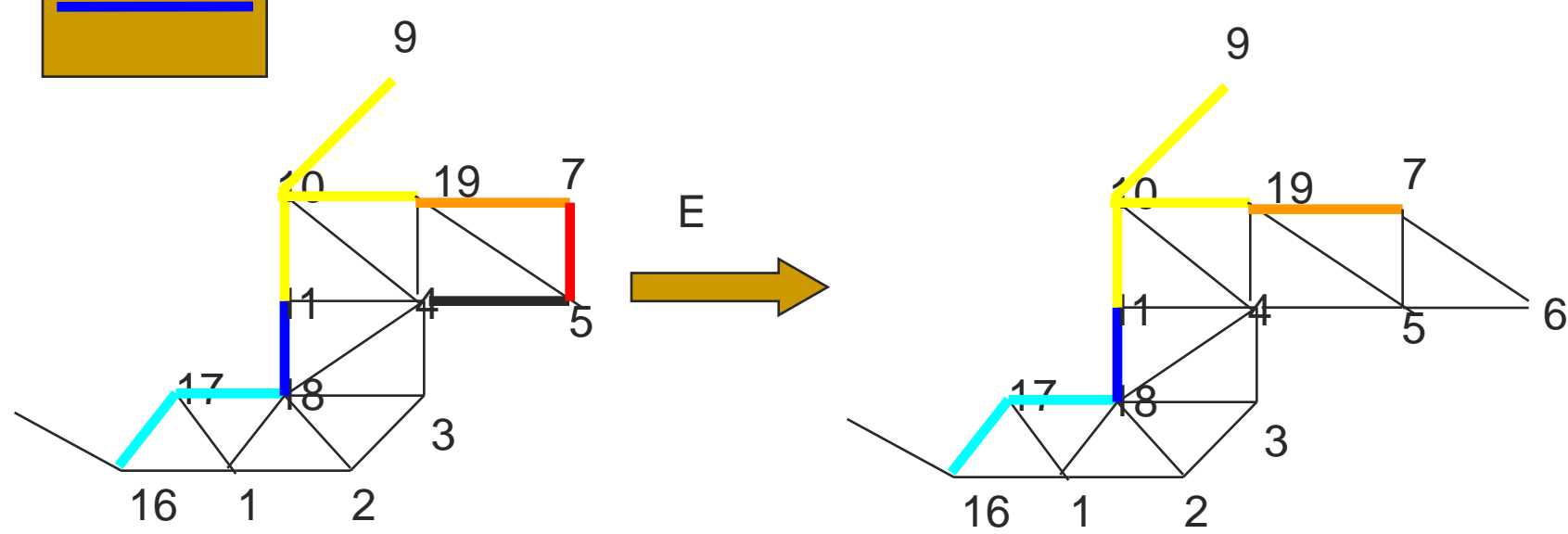
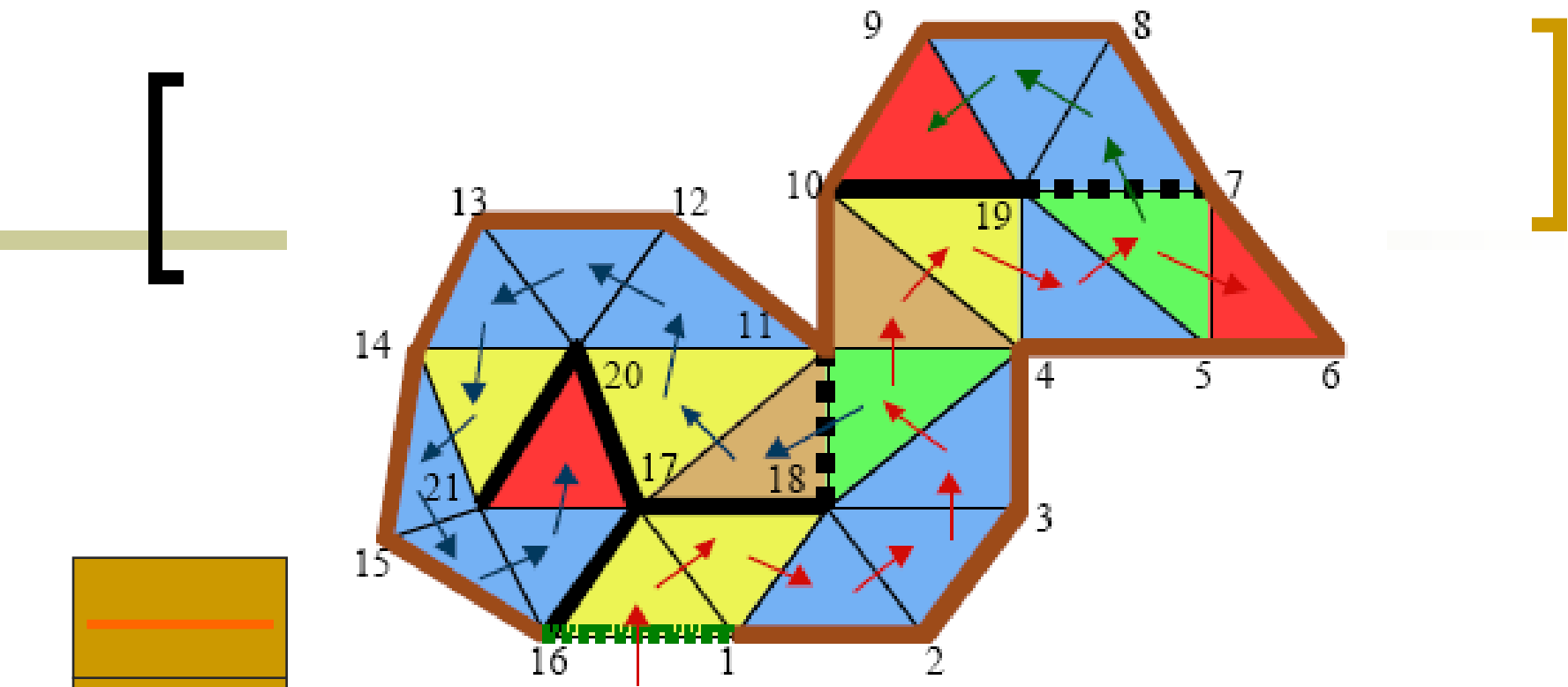














Face Fixer

[Face Fixer]

- Can handle arbitrary polygon meshes.
 - Quadrangular meshes can be compressed more efficient than their triangulated counterparts.
- Take properties and structures into account.

[Face Fixer]

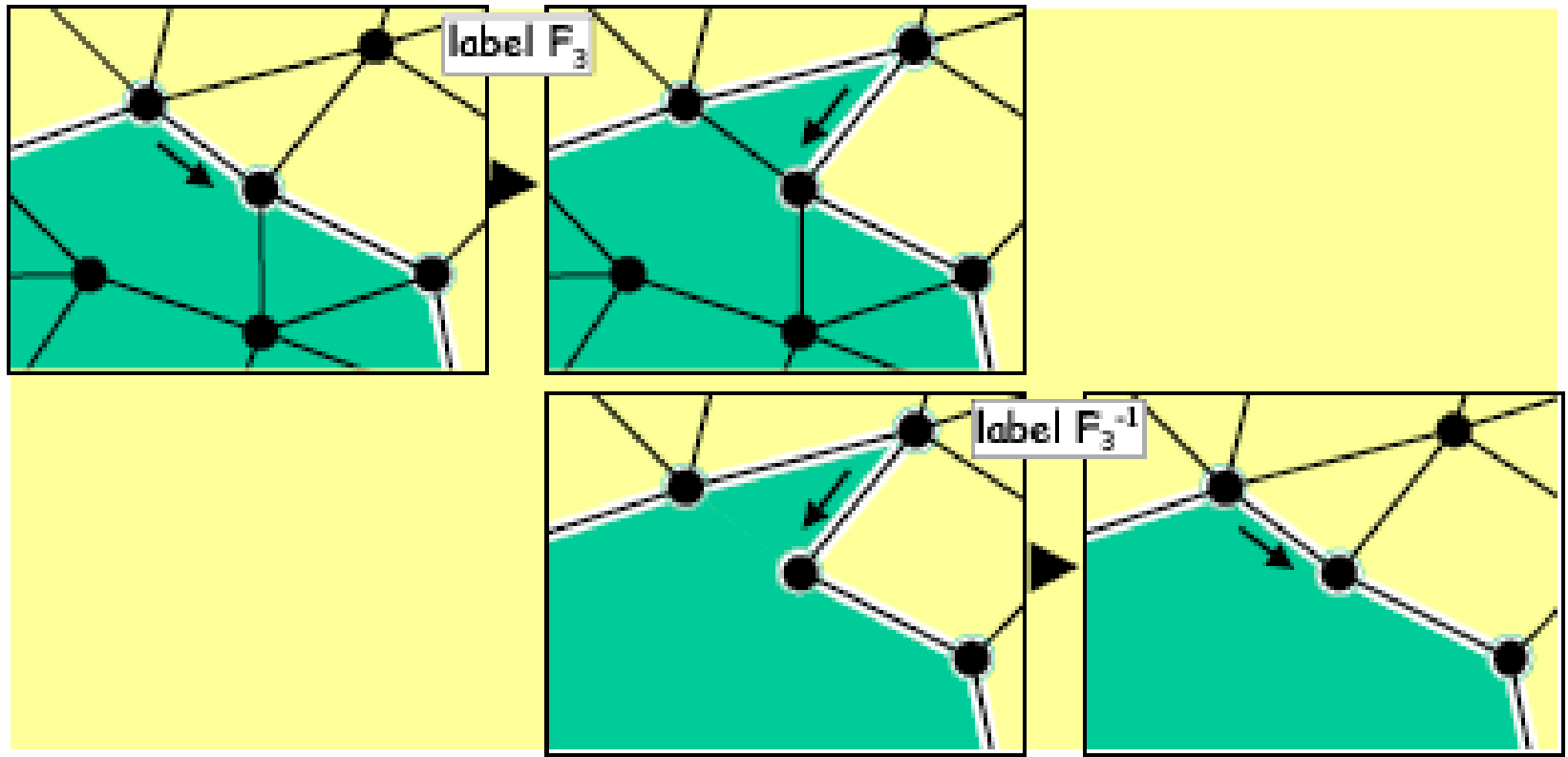
- Instead of take faces off, Face Fixer method take edge off.
- The total number of operations equal to total number of edges.

[Compression]

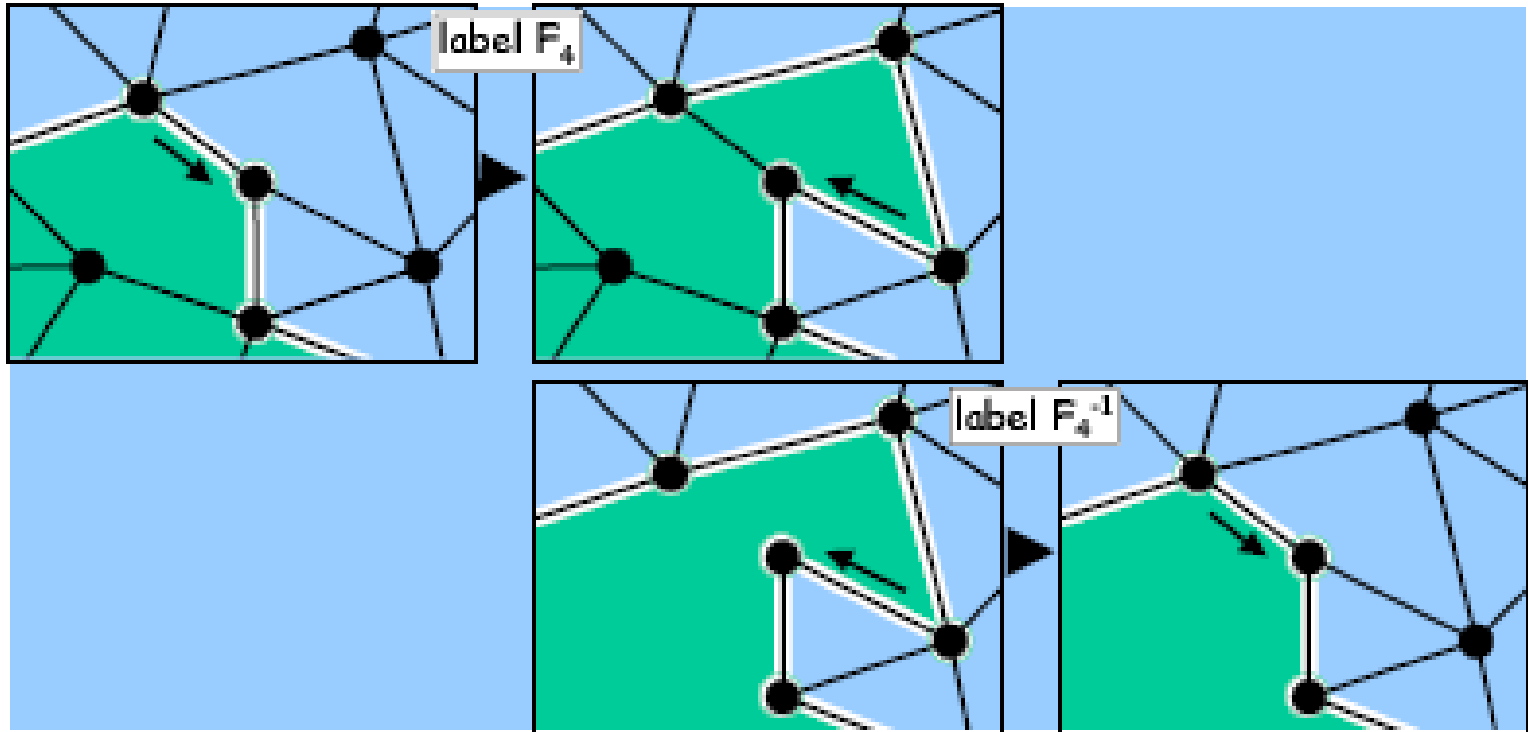
- Variable

- f : face count (0)
- h : hole count (0)
- v : vertex count (0)
- e : edge count (0)
- h : handle count (0)

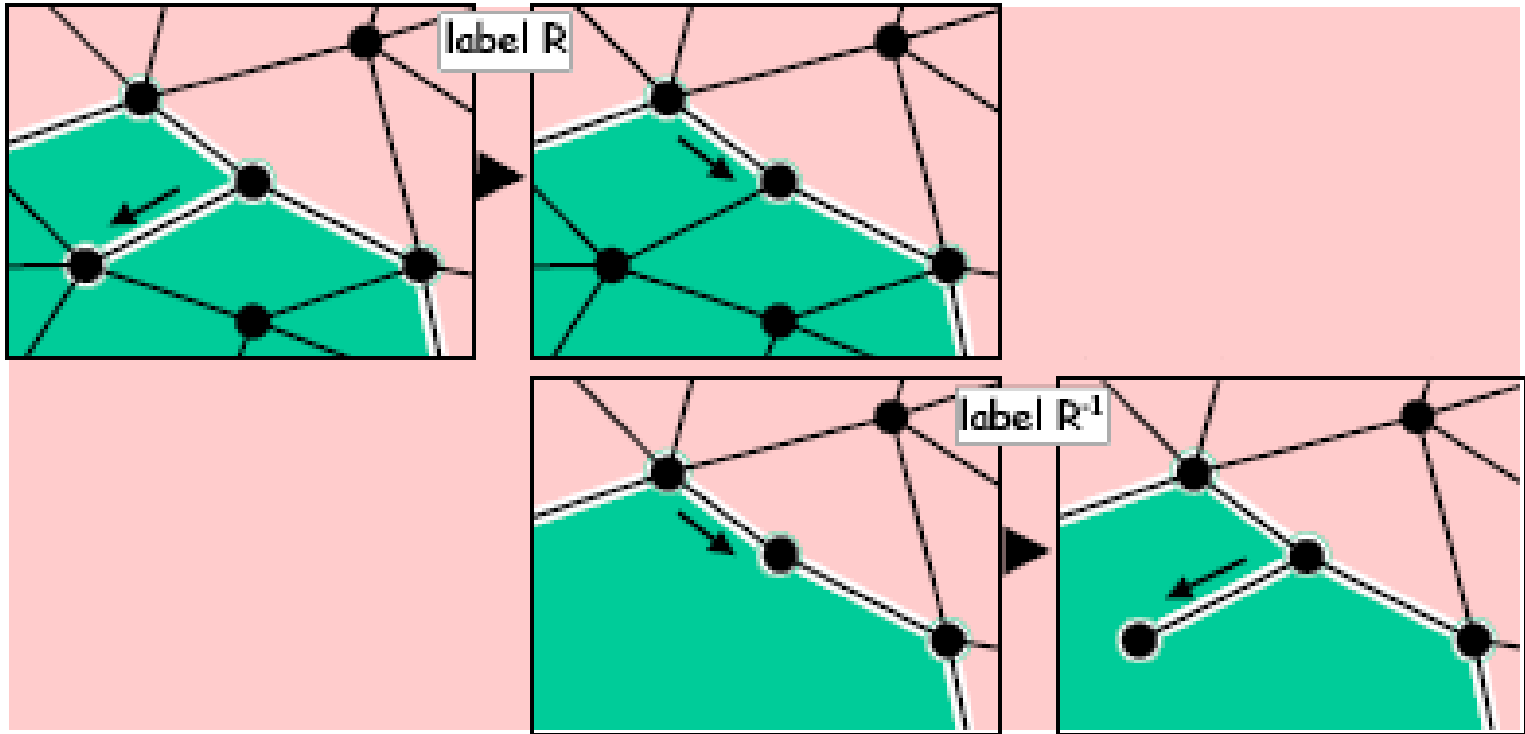
[Case F3]



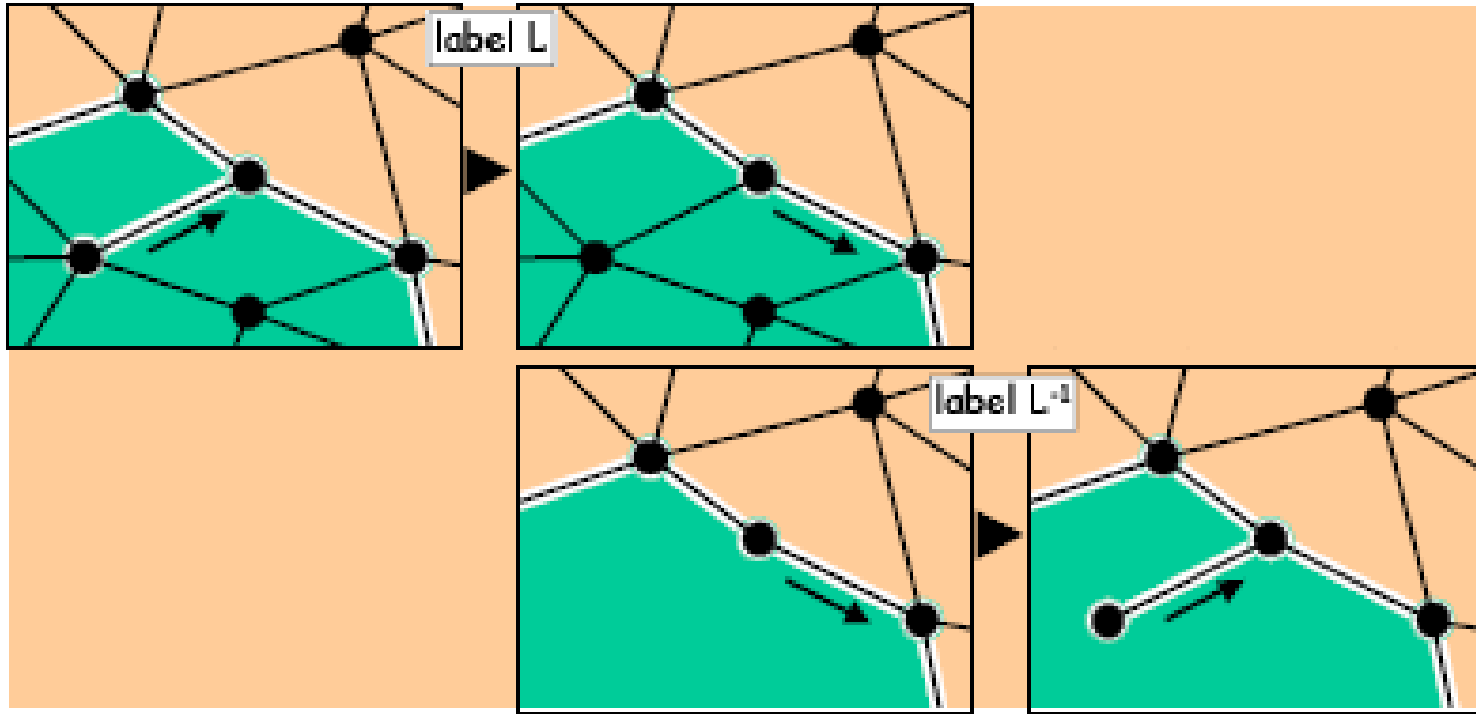
[Case F4]



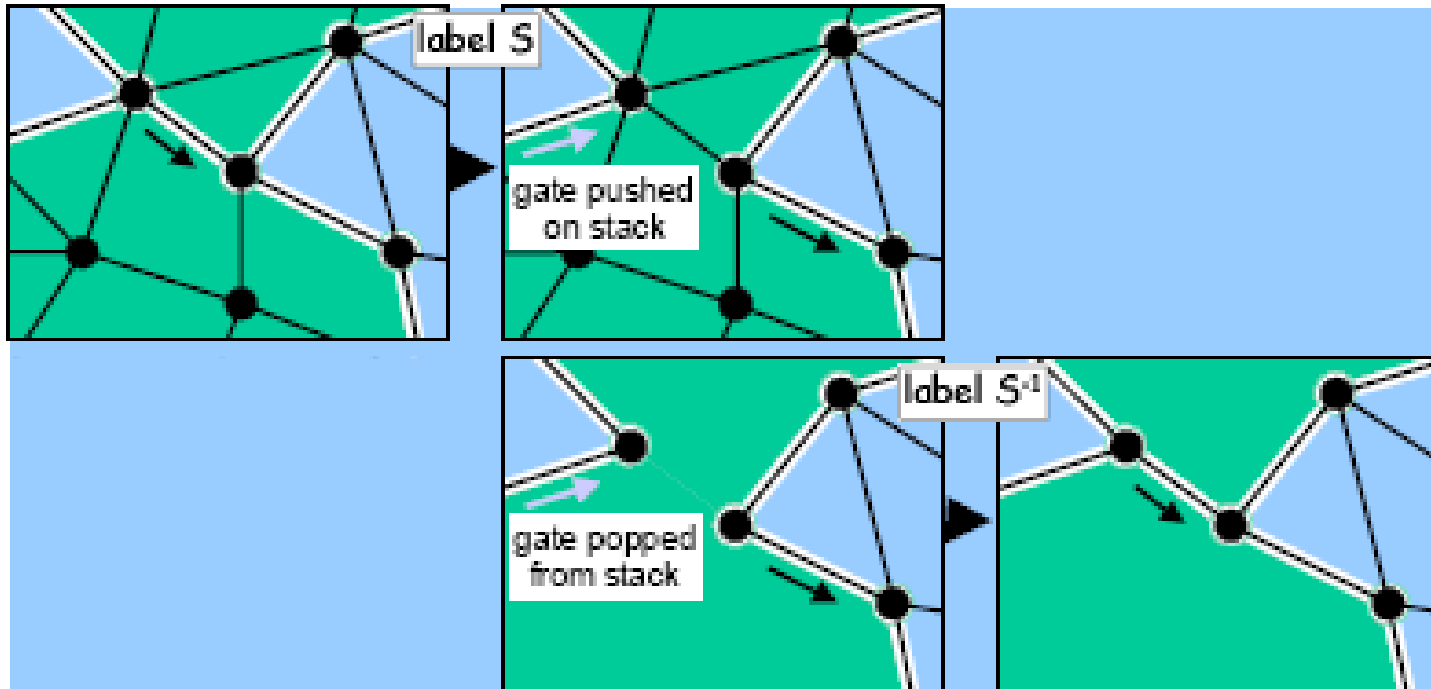
[Case R]



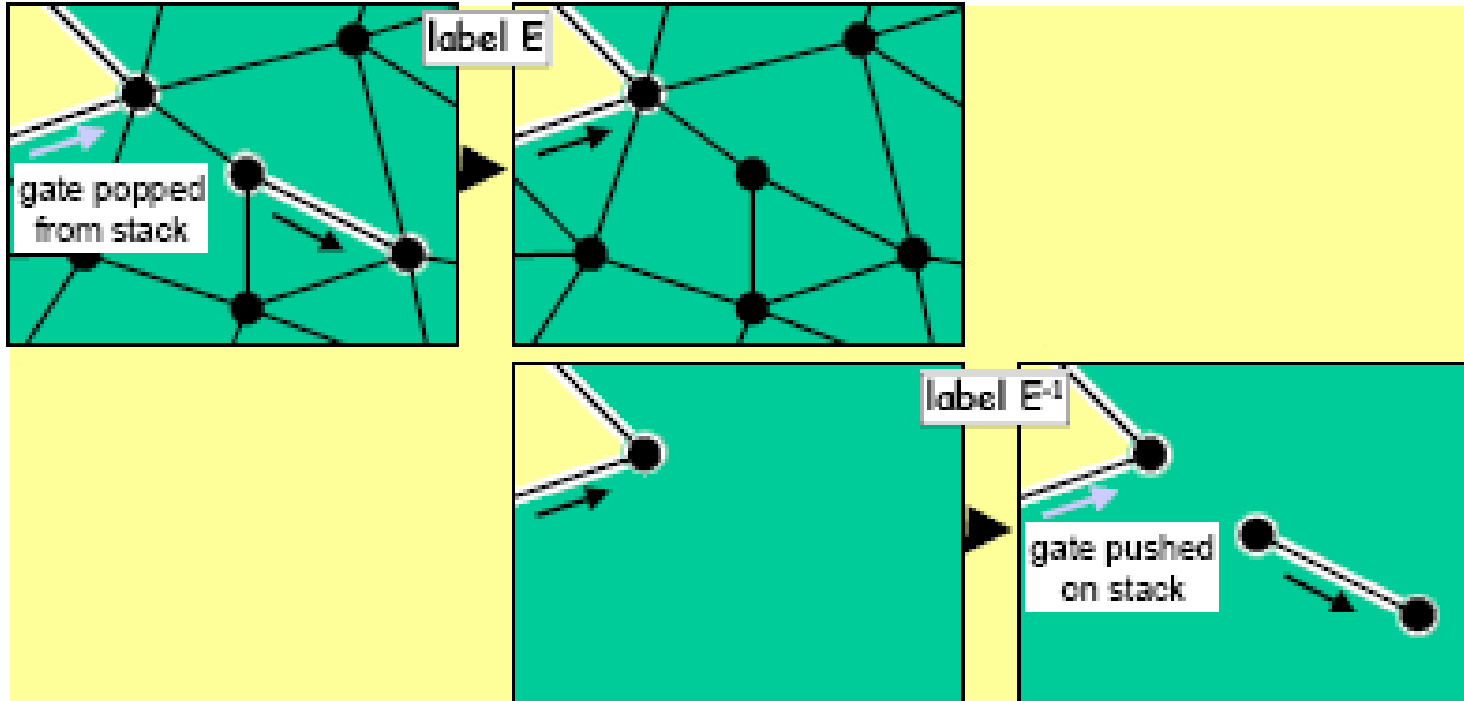
[Case L]



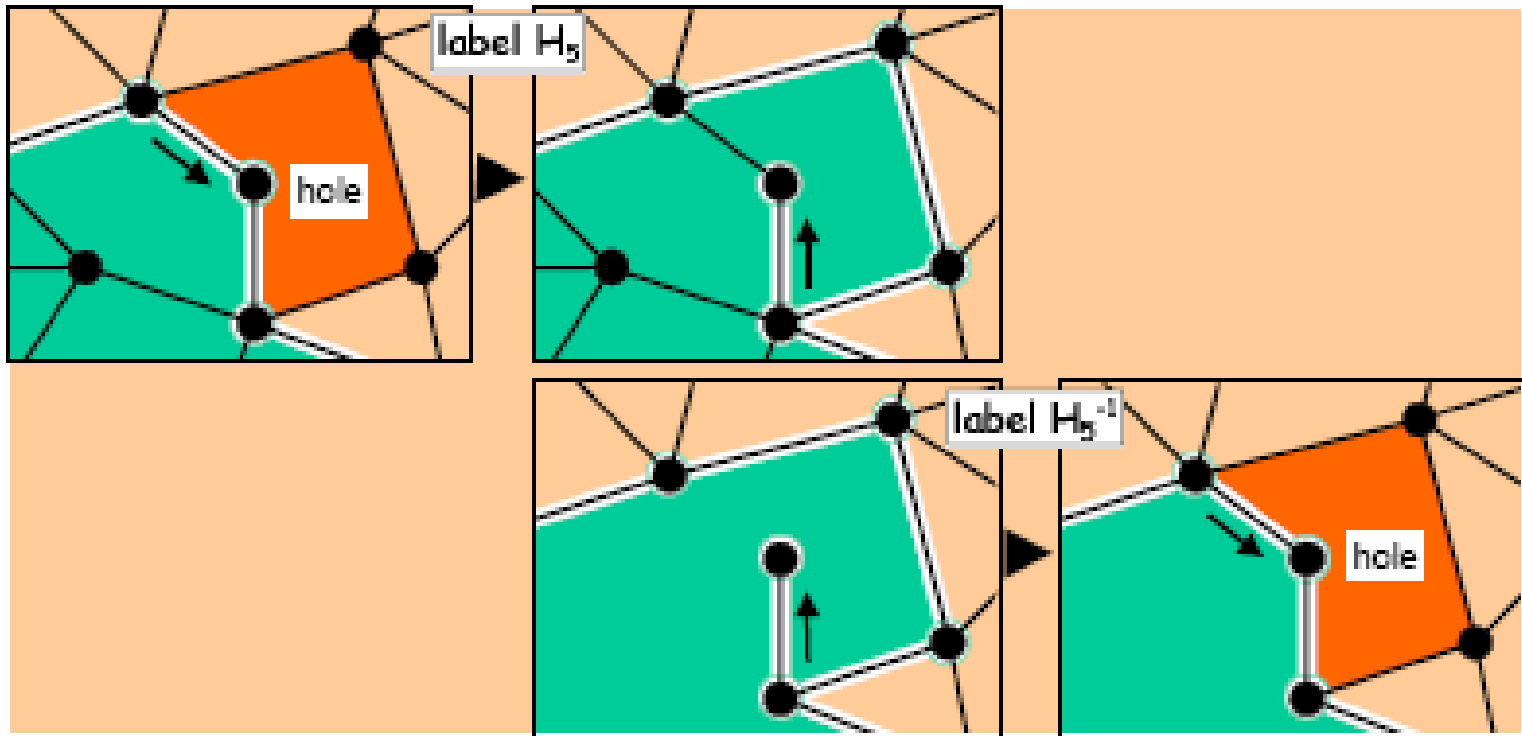
[Case S]



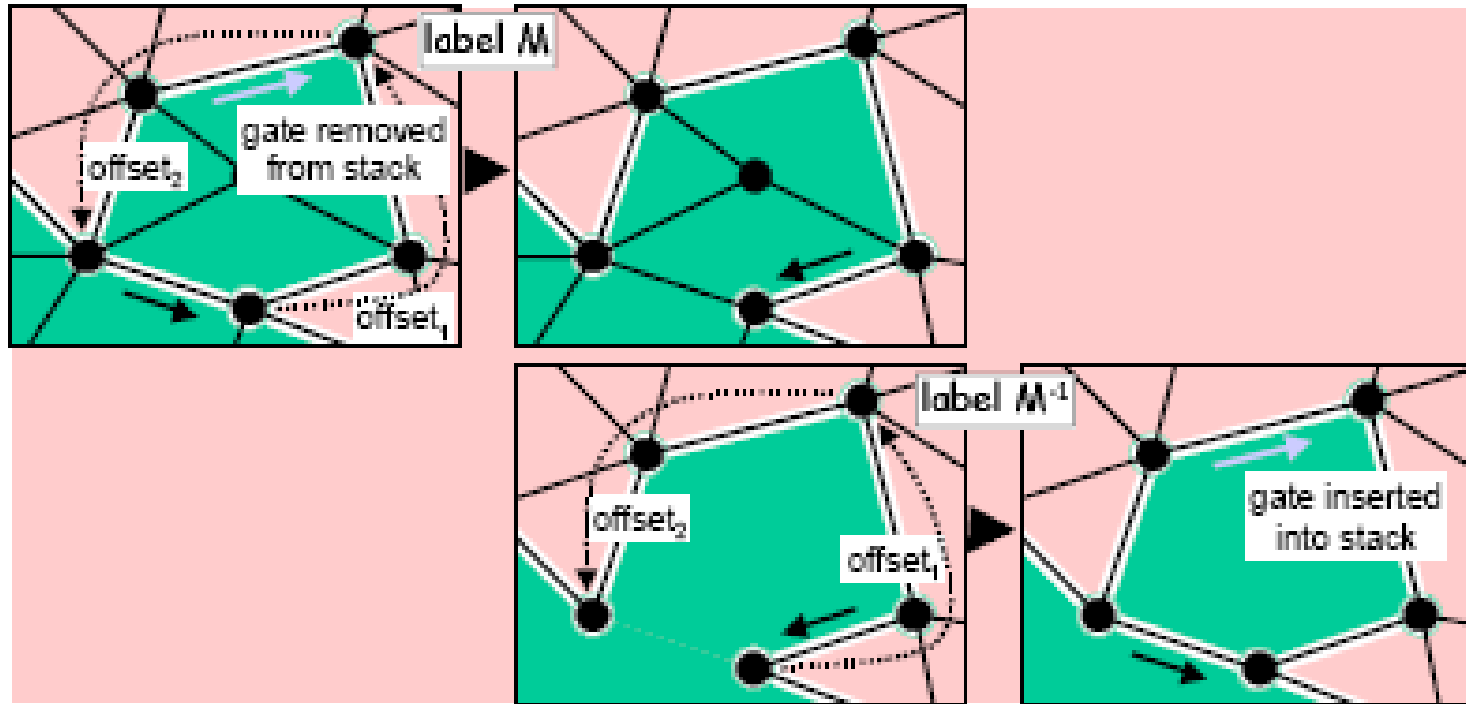
[Case E]



Case H5



[Case M]



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